

Convex Sets

Reference: Chapter 2, Boyd & Vandenberghe

Why convex?

$$x^* = \underset{g(x) \leq 0}{\operatorname{arg\,min}} f(x)$$

- Linear programming: f & g are both linear
- Quadratic programming
- Semidefinite programming
- Convex optimization: f is convex
constraints form a convex set

Lines, and line segments

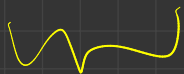
$$\underline{x}_1, \underline{x}_2 \in \mathbb{R}^n$$

$$\{ \alpha \underline{x}_1 + (1-\alpha) \underline{x}_2 \mid \alpha \in \mathbb{R} \}$$

is the line
passing through
 $\underline{x}_1, \underline{x}_2$

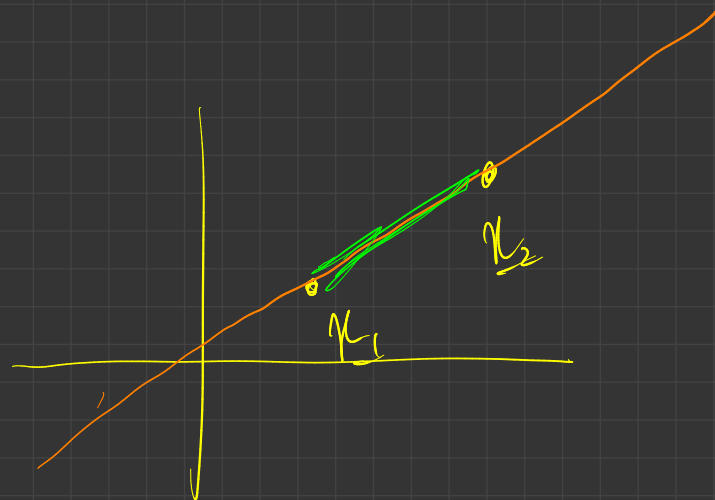


$$\underline{x}_1 + (1-\alpha) (\underline{x}_2 - \underline{x}_1)$$



direction

base



$$\frac{y - y_1}{x - x_1} = \frac{y_2 - y_1}{x_2 - x_1} \quad \begin{bmatrix} x - x_1 \\ y - y_1 \end{bmatrix} = \begin{bmatrix} \alpha(x_2 - x_1) \\ \alpha(y_2 - y_1) \end{bmatrix}$$

$$\begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} x_1 \\ y_1 \end{bmatrix} + \alpha \begin{bmatrix} x_2 - x_1 \\ y_2 - y_1 \end{bmatrix}$$

$$= \begin{bmatrix} x_1 \\ y_1 \end{bmatrix} + \alpha \begin{bmatrix} x_2 \\ y_2 \end{bmatrix} - \alpha \begin{bmatrix} x_1 \\ y_1 \end{bmatrix}$$

$$= \alpha \begin{bmatrix} x_2 \\ y_2 \end{bmatrix} + (1 - \alpha) \begin{bmatrix} x_1 \\ y_1 \end{bmatrix}$$

$$= \alpha \underline{x_2} + (1 - \alpha) \underline{x_1}$$

$$\{ \underline{x} = \alpha \underline{x}_1 + (1-\alpha) \underline{x}_2, \quad 0 \leq \alpha \leq 1 \}$$

is the line segment joining \underline{x}_1 & \underline{x}_2

Affine sets

A set S is affine if

$$\underline{x}_1, \underline{x}_2 \in S,$$

then

$$\alpha \underline{x}_1 + (1-\alpha) \underline{x}_2 \in S$$

$$\forall \alpha \in \mathbb{R}.$$

Consider \mathbb{R}^2 .

— Any straight line is affine

— $\{ \underline{x} \}$ is affine

— \mathbb{R}^2

— $A\underline{x} = \underline{b}$

$\underline{x}_1, \underline{x}_2$ are solutions of $A\underline{x} = \underline{b}$

$$\begin{aligned} A(\alpha \underline{x}_1 + (1-\alpha) \underline{x}_2) &= \alpha A\underline{x}_1 + (1-\alpha) A\underline{x}_2 \\ &= \alpha \underline{b} + (1-\alpha) \underline{b} \\ &= \underline{b} \end{aligned}$$

Affine combination of $\underline{x}_1, \underline{x}_2, \dots, \underline{x}_k$

$$\alpha_1 \underline{x}_1 + \alpha_2 \underline{x}_2 + \dots + \alpha_k \underline{x}_k$$

$$\sum_{i=1}^k \alpha_i = 1$$

Every affine set is a shift of a vector subspace

Consider S affine - $\underline{c} \in S$

Claim: $S - \underline{c} = \{ \underline{x} - \underline{c} : \underline{x} \in S \}$ is a vector subspace

$$\begin{aligned} \alpha (\underline{x}_1 - \underline{c}) + \beta (\underline{x}_2 - \underline{c}) &= \alpha \underline{x}_1 + \beta \underline{x}_2 - (\alpha + \beta) \underline{c} \\ &= \underbrace{\alpha \underline{x}_1 + \beta \underline{x}_2 + (1 - \alpha - \beta) \underline{c}}_{\text{Affine comb of } \underline{x}_1, \underline{x}_2, \underline{c} \text{ in } S} - \underline{c} \end{aligned}$$

Affine comb of $\underline{x}_1, \underline{x}_2, \underline{c}$ in S & hence

Dimension of affine space

$$\dim(S) = \dim(S - \underline{c})$$

↳ vector space dimension

Every affine set is the solution space of a system of linear equations

Every affine set can be written as

$$S = \underline{v} + \underline{b}$$

$$\downarrow$$
$$\{ \underline{x} : A\underline{x} = \underline{0} \}$$

$$S = \{ \underline{x} + \underline{b} : A\underline{x} = \underline{0} \}$$

$$= \{ \underline{y} : A(\underline{y} - \underline{b}) = \underline{0} \}$$

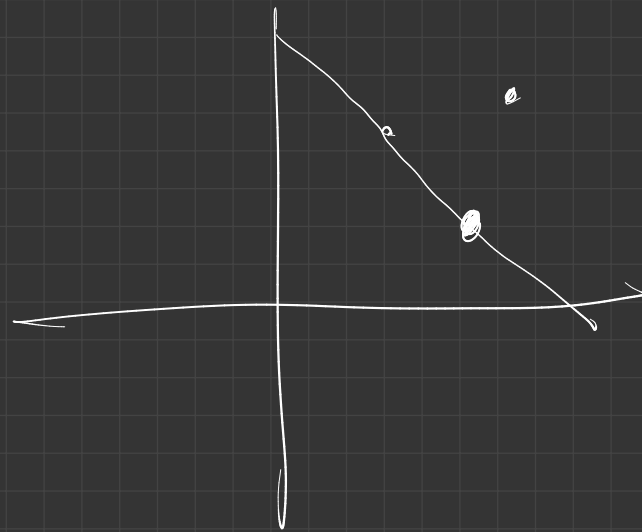
$$= \{ \underline{y} : A\underline{y} = A\underline{b} \}$$

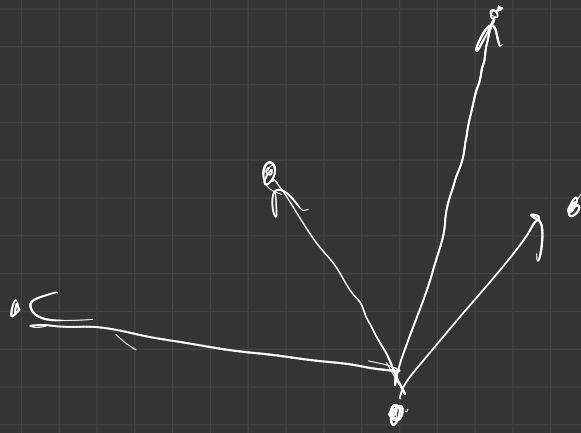
Affine hull, examples

Given a set $C \subseteq \mathbb{R}^n$, the affine hull

$$\text{aff}(C) = \left\{ \underline{x} = \sum_{i=1}^k \alpha_i \underline{x}_i \mid \underline{x}_1, \dots, \underline{x}_k \in \mathbb{R}^n, \sum_{i=1}^k \alpha_i = 1 \right\}$$

→ Smallest affine set that contains C





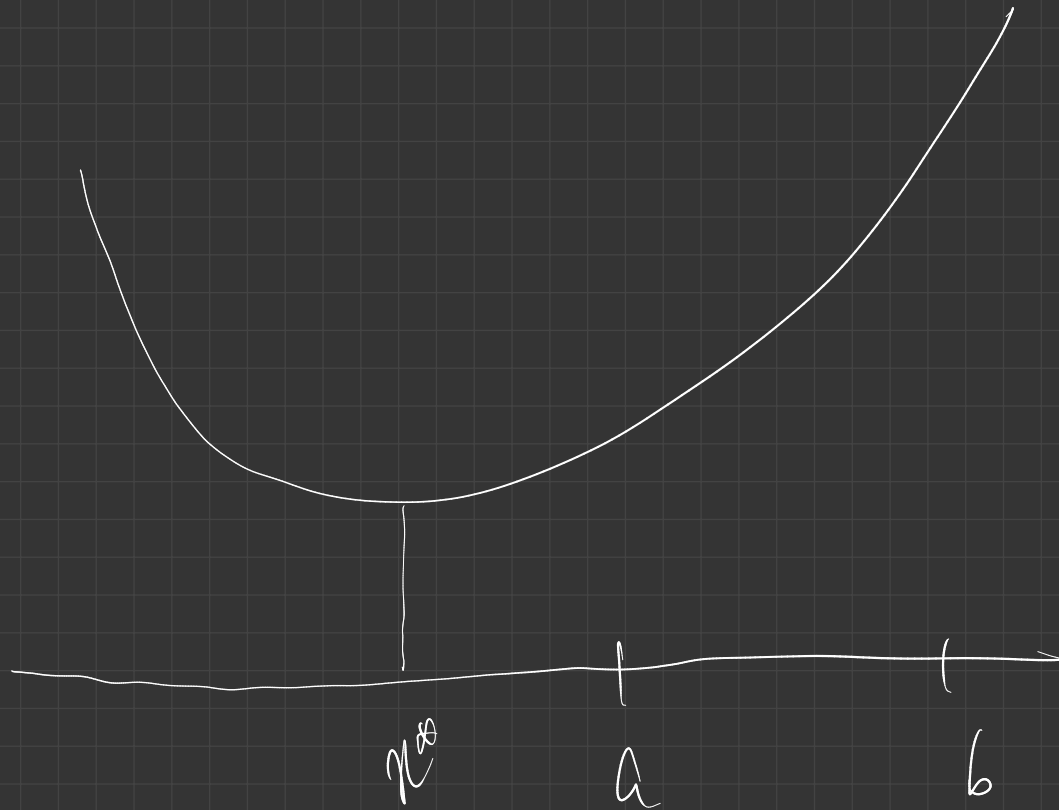
no

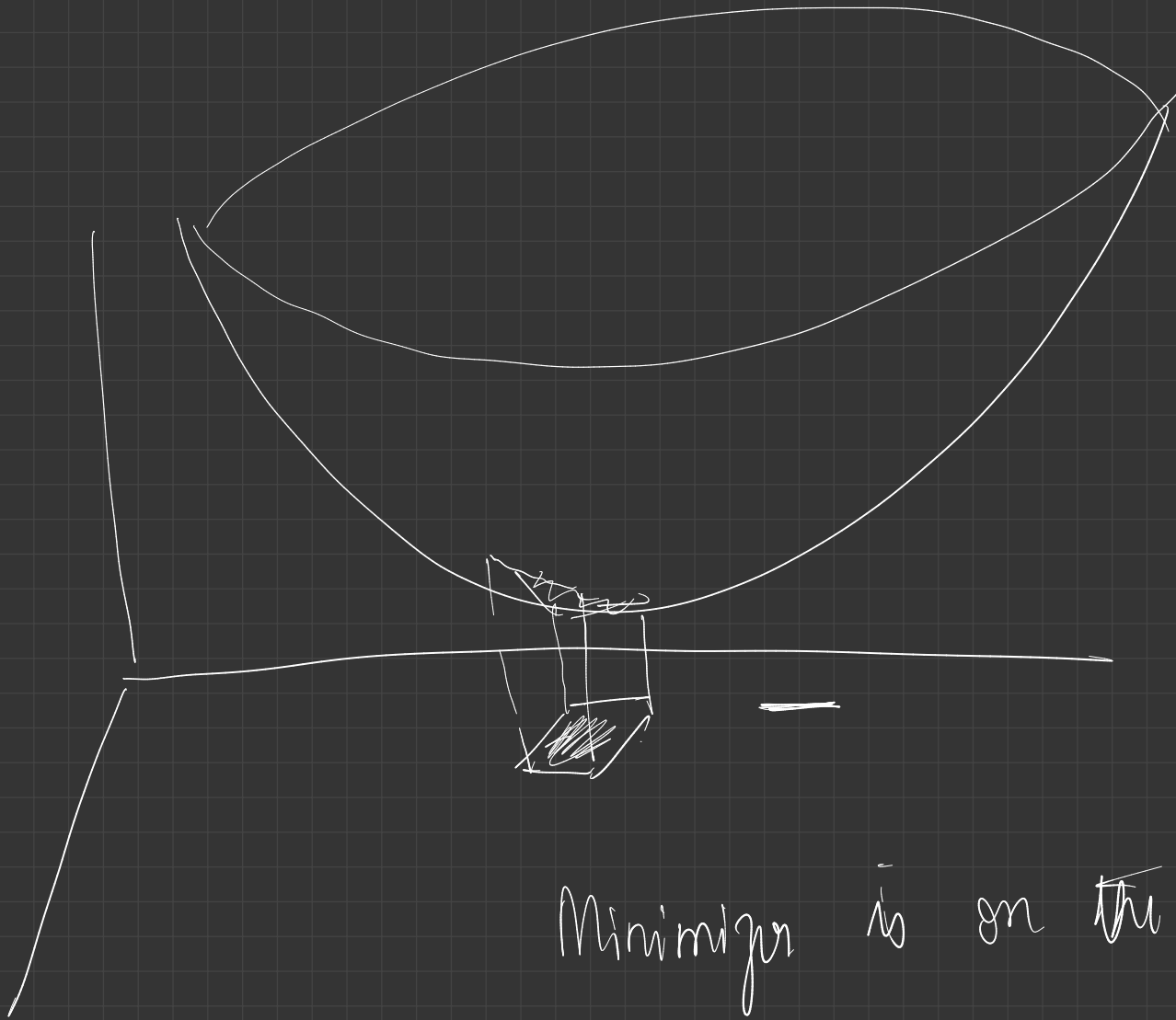
P I

argmin $f(x)$
 $x \in [a, b]$

P II

argmin $f(x)$

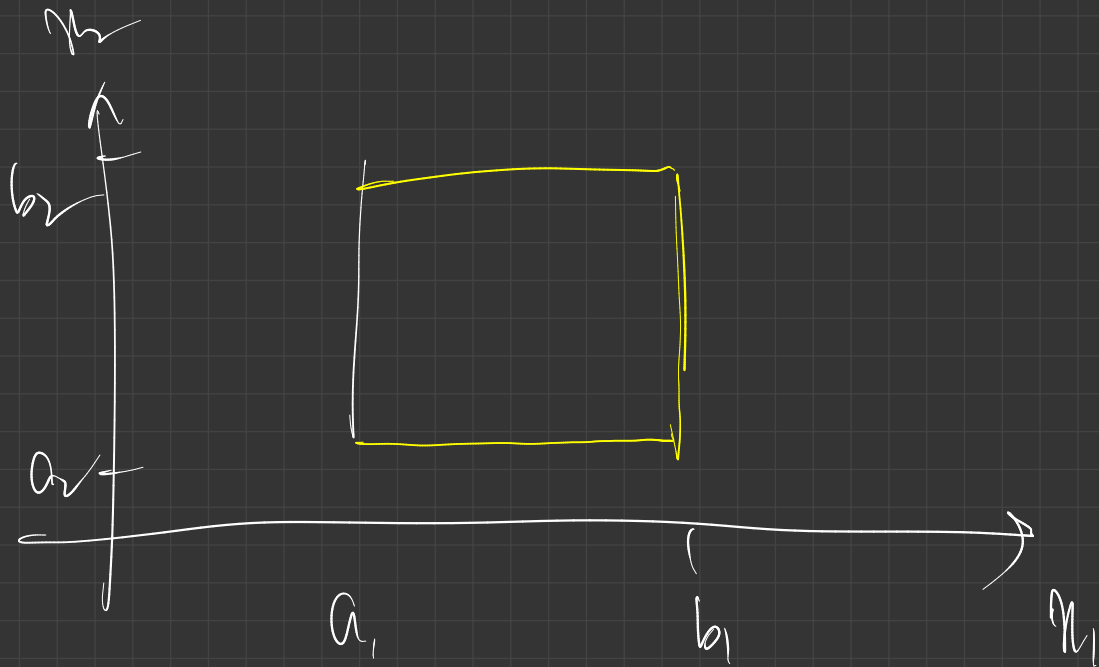




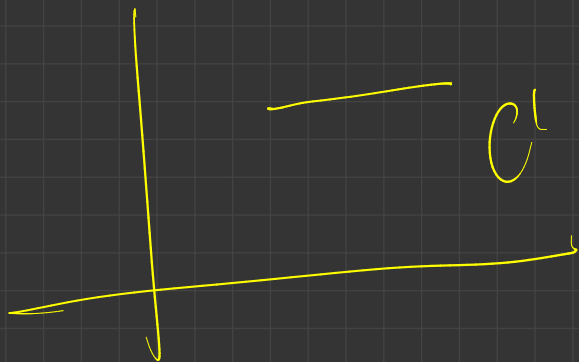
Minimizer is on the boundary

$$C = \left\{ \begin{array}{l} a_1 \leq x_1 \leq b_1 \\ a_2 \leq x_2 \leq b_2 \end{array} \right\}$$

$$\text{int}(C) = \left\{ \begin{array}{l} a_1 < x_1 < b_1 \\ a_2 < x_2 < b_2 \end{array} \right\}$$



$bd(C)$

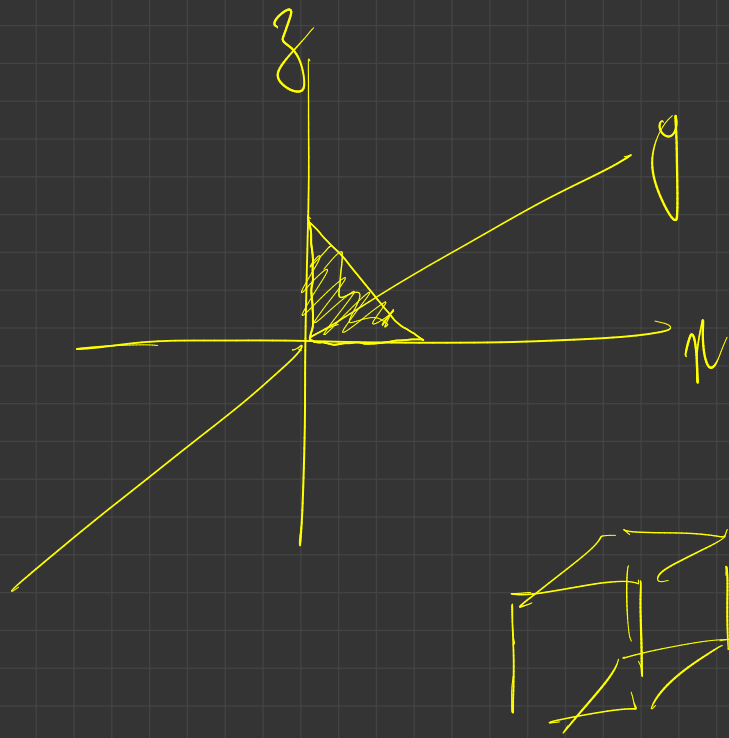


$$\text{int}(C') = \emptyset$$

Relative interior and boundary of a set

$$\text{Relint}(C) = \left\{ \underline{x} \mid \begin{array}{l} B_r(\underline{x}, \varepsilon) \cap \text{aff}(C) \subseteq C \\ \text{for some } \varepsilon > 0 \end{array} \right\}$$

$$\text{bd}(C) = \text{closure}(C) \setminus \text{relint}(C)$$



Convex combinations, convex sets and convex hulls

$$\alpha_1 \underline{x}_1 + \alpha_2 \underline{x}_2 + \dots + \alpha_k \underline{x}_k \quad \text{with} \quad \sum_{i=1}^k \alpha_i = 1$$

Convex combination of $\underline{x}_1, \dots, \underline{x}_k$ $\alpha_i \geq 0 \quad \forall i$

C is convex if $\underline{x}_1, \underline{x}_2 \in C \Rightarrow \alpha \underline{x}_1 + (1-\alpha) \underline{x}_2 \in C$
 $\forall \alpha \in [0, 1]$

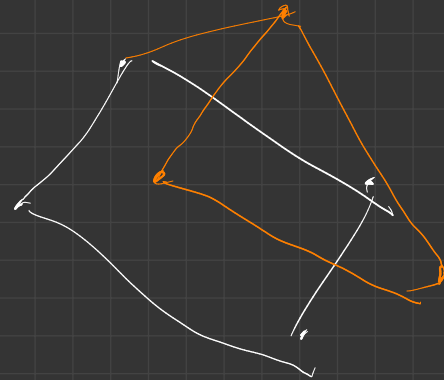
convex hull \vee $\text{conv}(C)$ = set of all convex comb. of pts in C

A, B

$$A \cap B = \emptyset$$

$$\text{conv}(A) \cup \text{conv}(B)$$

$$\text{conv}(A \cup B)$$



Infinite convex combinations

$$C \quad \underline{x}_1, \underline{x}_2, \dots, \underline{x}_n, \dots$$

$$\sum_{i \in I} \alpha_i \underline{x}_i \quad \text{with} \quad \alpha_i \geq 0 \quad \forall i$$

$\underbrace{\hspace{10em}}_{\substack{\text{if exists} \\ \sum \alpha_i = 1}}$

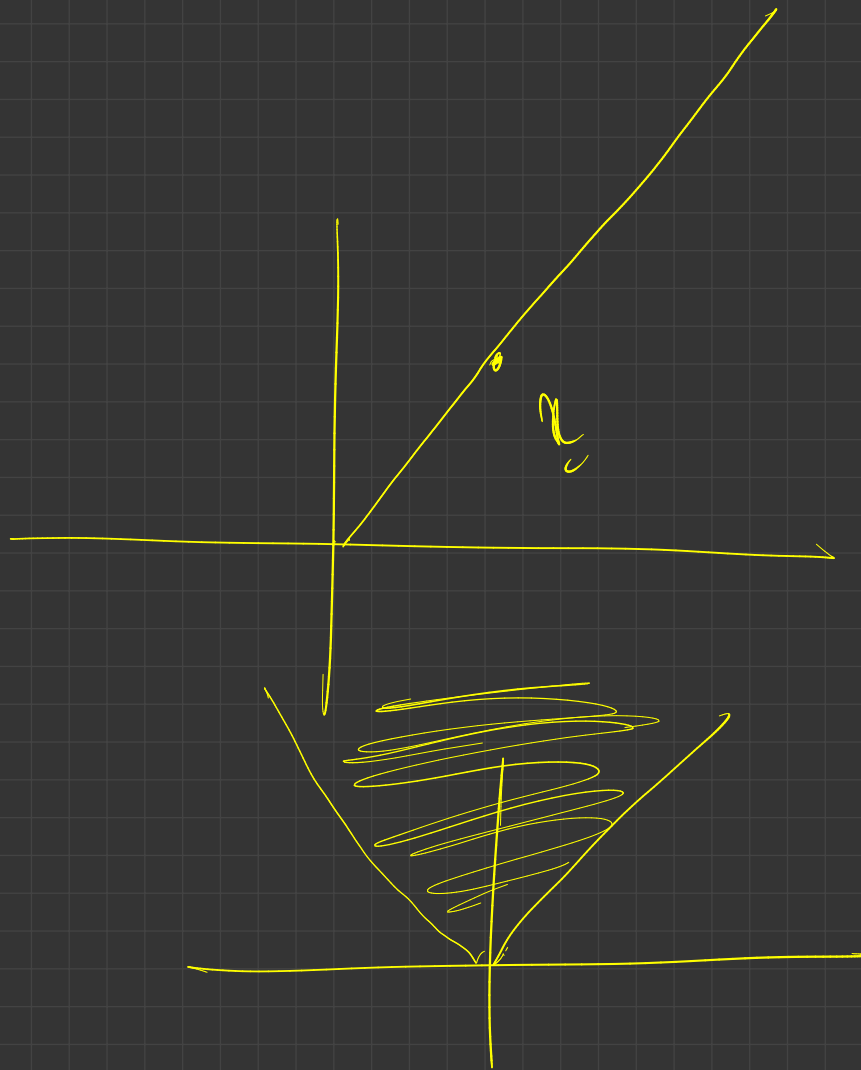
$$C \quad f \rightarrow \text{pdf on } C \quad \int_{\mathcal{N} \subset C} f(q) dq = 1$$

$$\int_{\mathcal{N} \subset C} \lambda f(q) dq$$

Cones and conic combinations

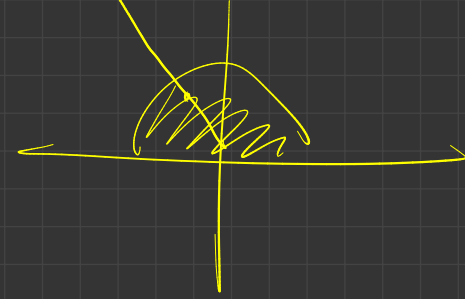
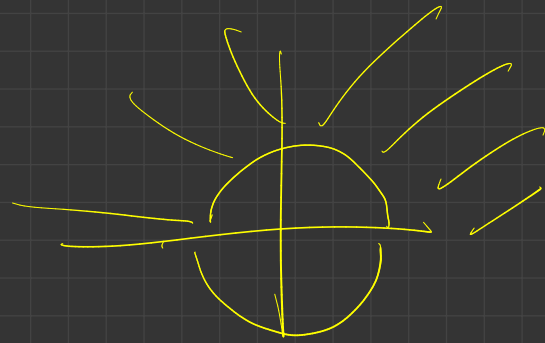
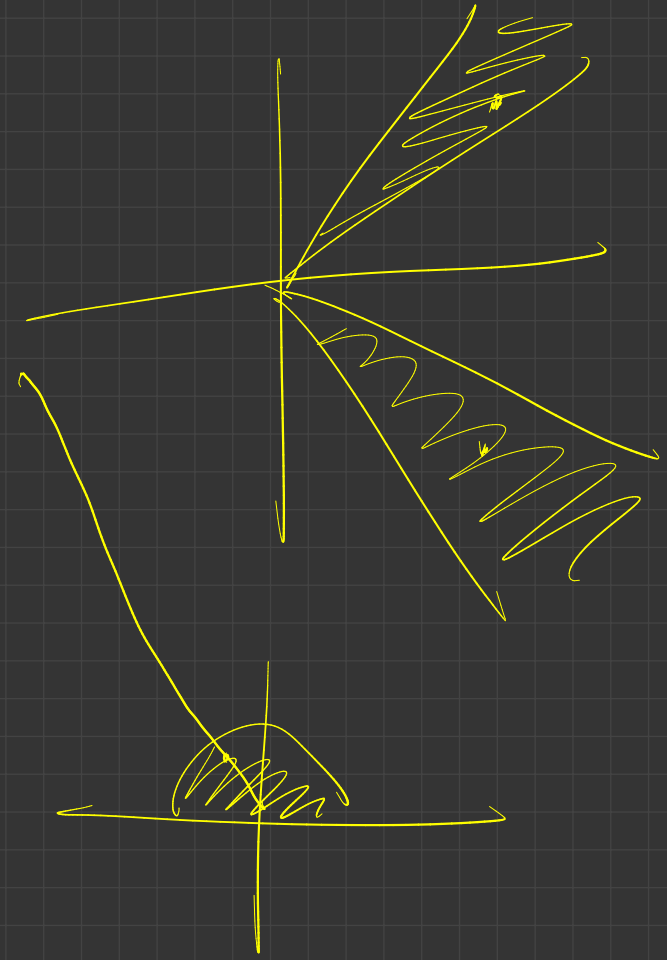
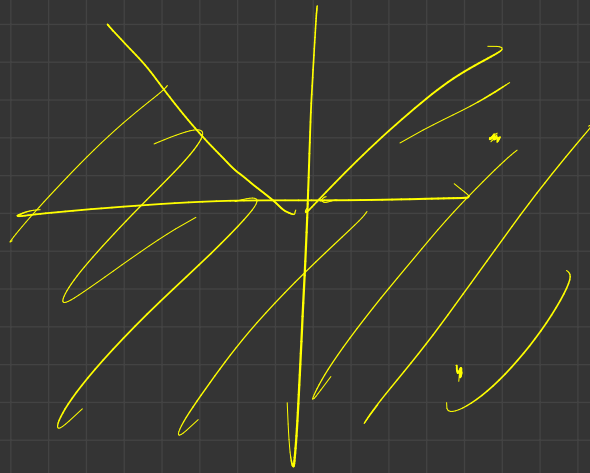
C is a cone if $x \in C$, then

$$\alpha x \in C \quad \forall \alpha \geq 0$$



$$f(x, u) = u \in \mathbb{R}$$

$$g(x, y) = y \geq |u|y$$

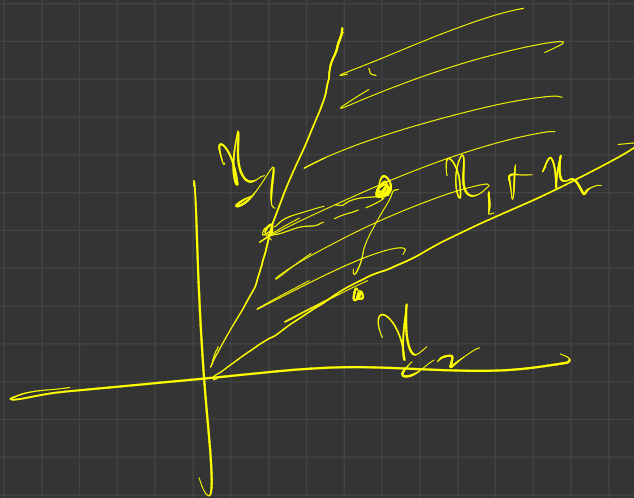


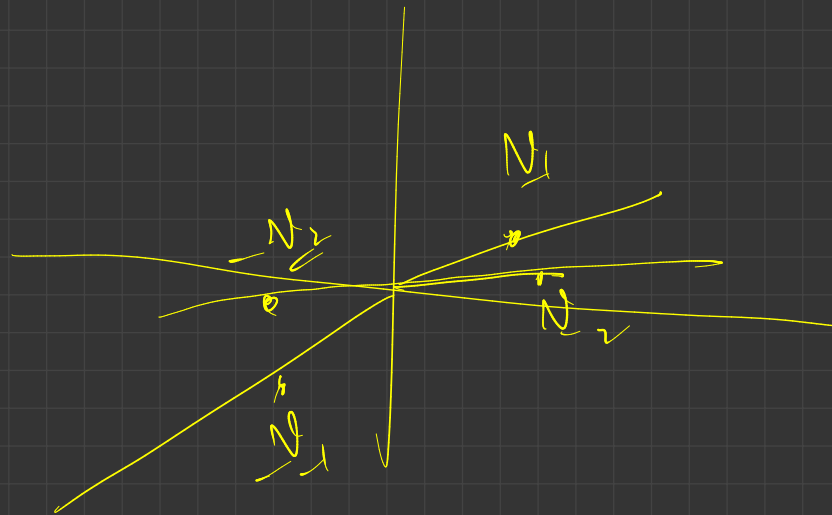
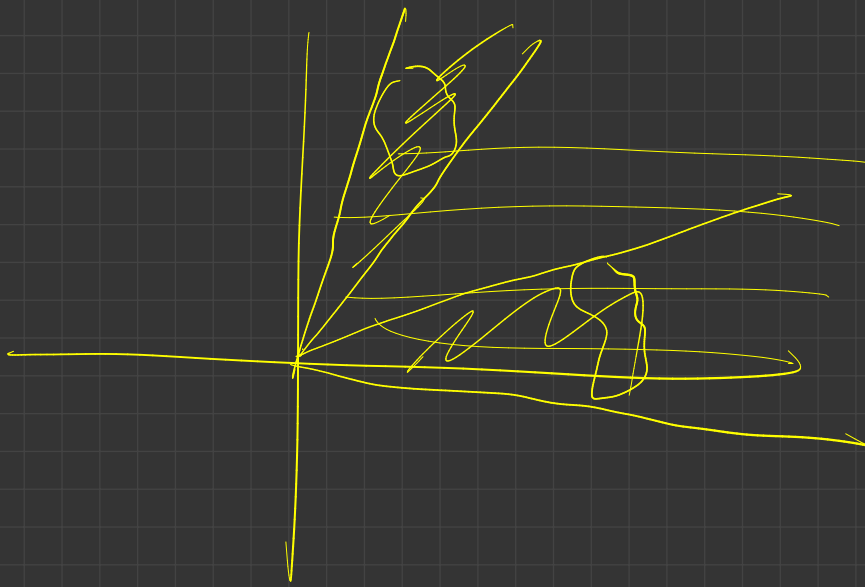
Conic hull

Conic combination:

$$r_1, r_2, \dots, r_k$$

$$\alpha_1 r_1 + \alpha_2 r_2 + \dots + \alpha_k r_k \quad \alpha_i \geq 0 \quad \forall i$$





$$\alpha_1 v_1 + \alpha_2 v_2$$

$$= \alpha_3 v_1 + \alpha_4 v_2$$

$$(\alpha_1 - \alpha_3) v_1 + (\alpha_2 - \alpha_4) v_2$$

Examples

convex

affine

line

① \emptyset

✓

✓

✓

② subspace of \mathbb{R}^n

✓

✓

✓

③ line segment

✓

X

X

④ $\{ \underline{x}_0 + \alpha(\underline{v} - \underline{x}_0) : \alpha \geq 0 \}$

✓

X

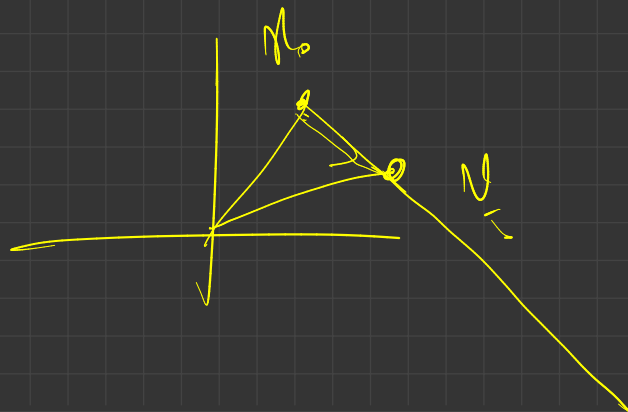
No general

$\underline{x}_0, \underline{v} \in \mathbb{R}^n$

Yes if

$\underline{x}_0 = \underline{v}$

Yes if $\underline{x}_0 = \underline{0}$



Hyperplanes and halfspaces

$$H = \{ \underline{x} : \underline{a}^T \underline{x} = b \}$$

$$\underline{a} \in \mathbb{R}^n, b \in \mathbb{R}$$



Affine set of dim $n-1$

$$\{ \underline{x} : \underline{a}^T \underline{x} \geq b \}$$

$$\{ \underline{x} : \underline{a}^T \underline{x} \leq b \}$$

Halfspace

$$a_1^T x \leq b_1$$

$$a_2^T x \leq b_2$$

!

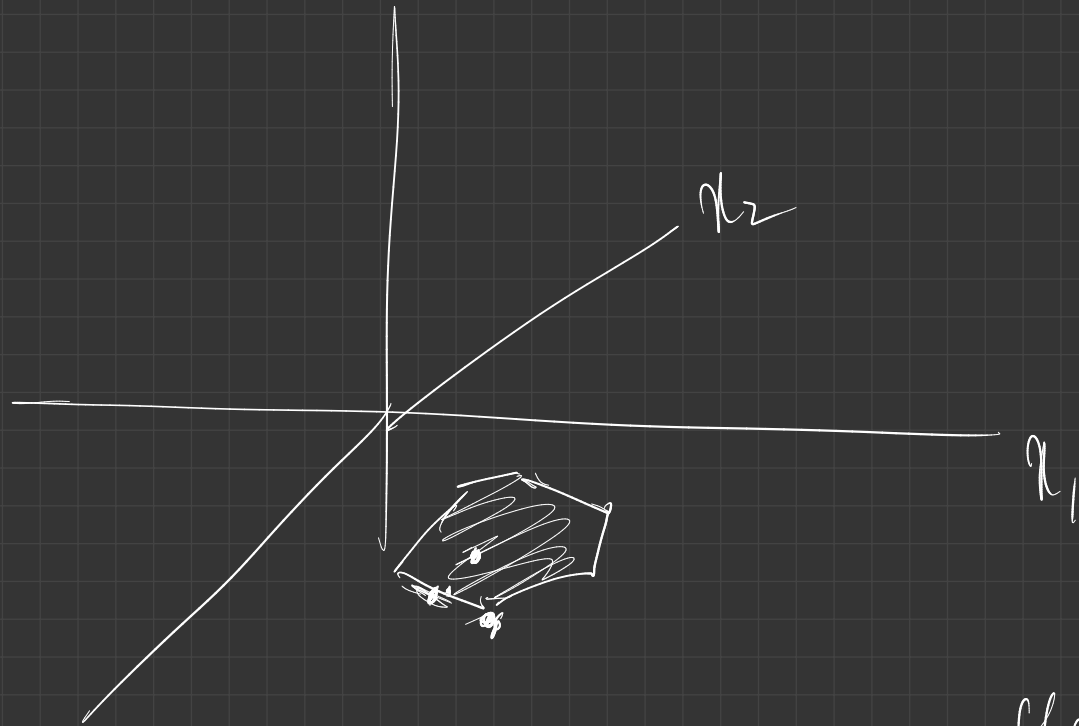
$$a_k^T x \leq b_k$$

$$a_i \in \mathbb{R}^n$$

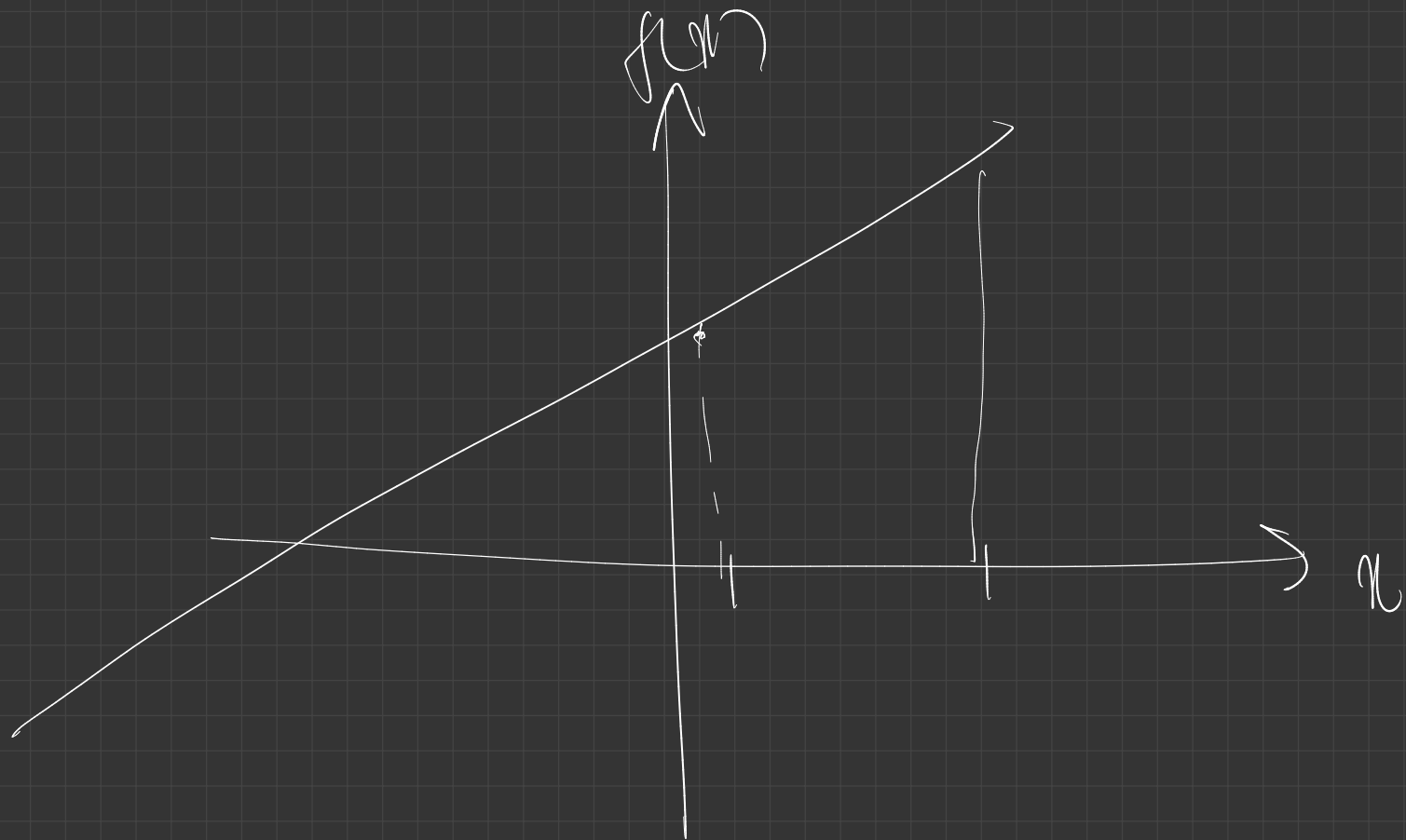
$$b_i \in \mathbb{R}$$



Polyhedron/polytope : Intersection of half spaces



$$f(x_1, x_2) = \alpha x_1 + \beta x_2$$



Linear program

$$f(x) = \underline{a}^T \underline{x}$$

$$\text{s.t.} \quad \underline{a}_1^T \underline{x} \leq b_1$$

$$\underline{a}_2^T \underline{x} \leq b_2$$

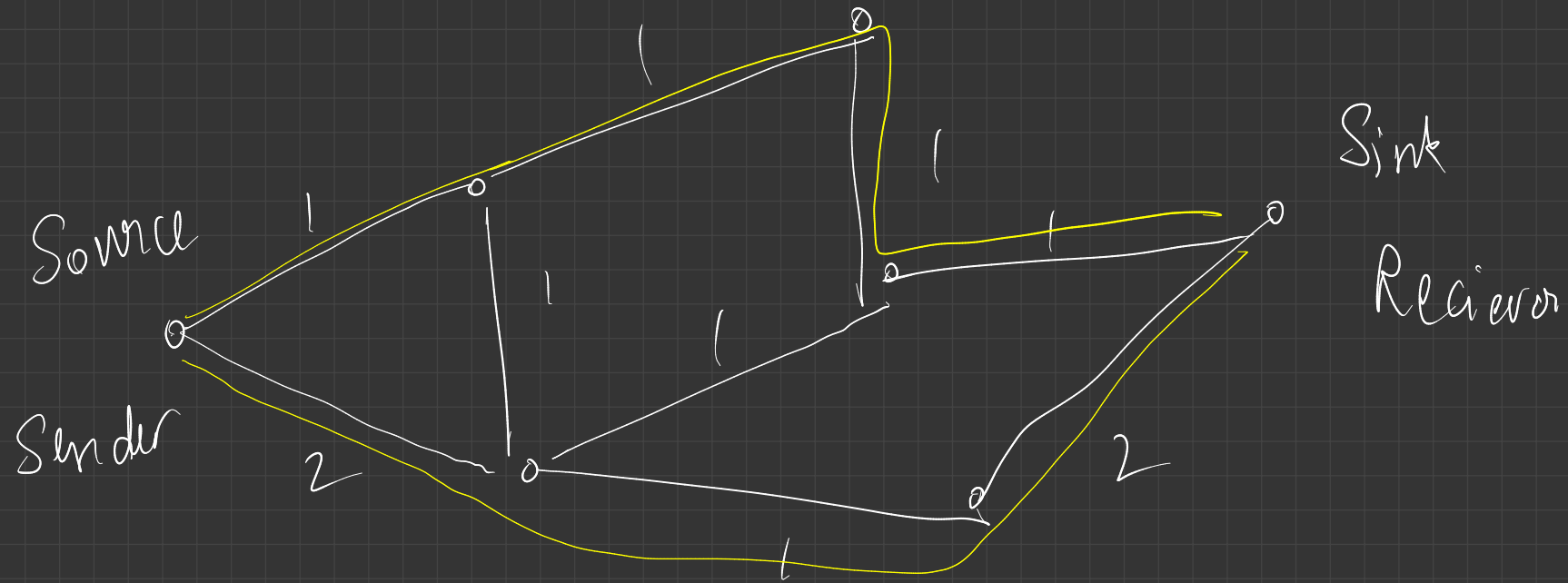
⋮

$$\underline{c}_1^T \underline{x} = d_1$$

$$\underline{c}_2^T \underline{x} = d_2$$

⋮

Max flow problem



Variables, Flows along each edge

(x_{ij}) : r_{ij} (# of packets sent from i to j per sec)

X - $n \times n$ matrix

Capacities, For each pair of vertices i, j

$$c_{ij} \geq 0$$

$$c_{ij} = c_{ji}$$

$$c_{ii} = 0$$

C - $n \times n$ matrix

Objective

$f(x)$

= Total flow leaving source / total flow into sink

$f(x)$

$$= \sum_{j=1}^{n-1} x_{0j} = \sum_{j=0}^{n-1} r_{j, n-1}$$

Constraints :

① Capacity constraint : $x_{ji} + x_{ij} \leq C_{ij}$

② $x_{ij} \geq 0 \quad \forall i, j$

③ $\forall i, \sum_{\substack{j=0 \\ j \neq i}}^{n-1} x_{ji} - \sum_{\substack{j=0 \\ j \neq i}}^{n-1} x_{ij} = 0$
 $1 \leq i \leq n-2$

④ $\sum_{j=1}^{n-1} x_{j0} = 0$

$\sum_{j=0}^{n-2} x_{n-1,j} = 0$

Variable: X : $n \times n$ matrix

Objective: $\text{Max} \{ \text{sum}(X[0, :]) \}$

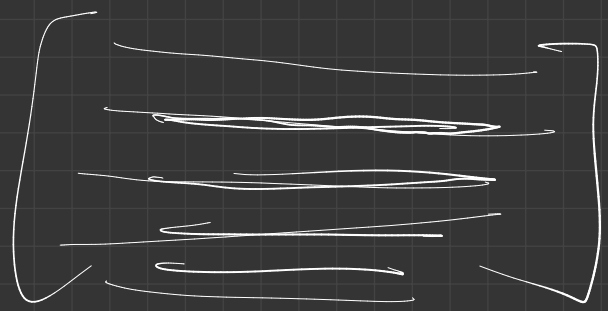
Constraints: $X \geq 0$ (elementwise)

$$X + X^T \leq C$$

$$\text{sum}(X[1:n-2, :], \text{axis}=1) - \text{sum}(X^T[1:n-2, :], \text{axis}=1) = 0$$

total inflow
↑

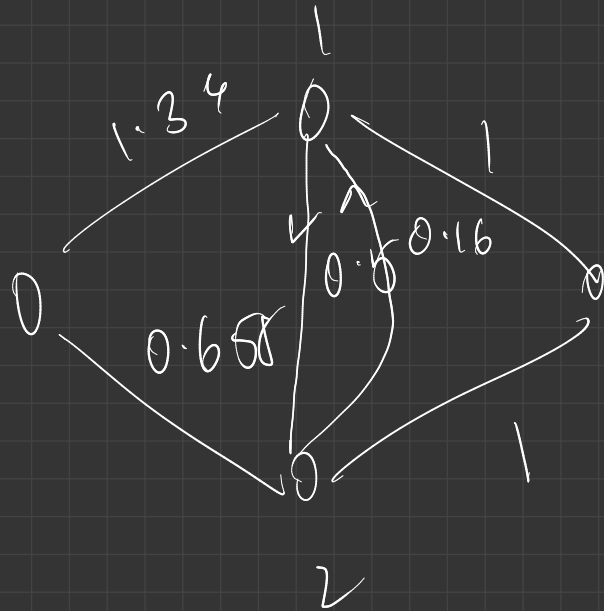
total outflow
from a node
↙



$$\text{sum}(X[:, 0]) = 0$$

$$\text{sum}(X[n-1, :]) = 0$$

$$X^{\infty} = \begin{bmatrix} 0 & 1.34 & 0.658 & 0 \\ 0 & 0 & 0.5 & 1 \\ 0 & 0.16 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$



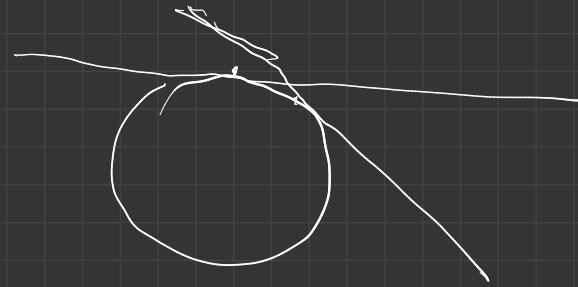
$$f(x, y) = x^2 + y^2, \text{ Minimiere.}$$

$$-1 \leq x \leq 1$$

$$-1 \leq y \leq 1$$

Two different ways of looking at a closed convex set

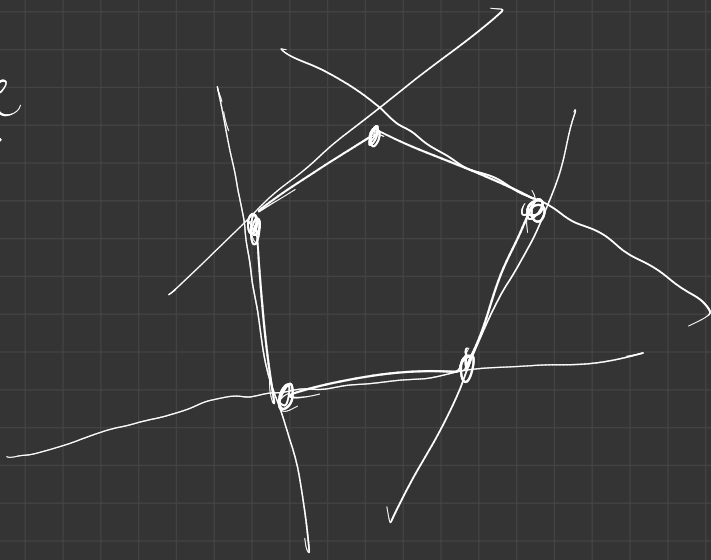
$C = \bigcap \{H \mid H \text{ a halfspace that contains } C\}$



→ Halfspace description
of a closed
convex set

$C = \text{conv}(C')$ → convex hull description

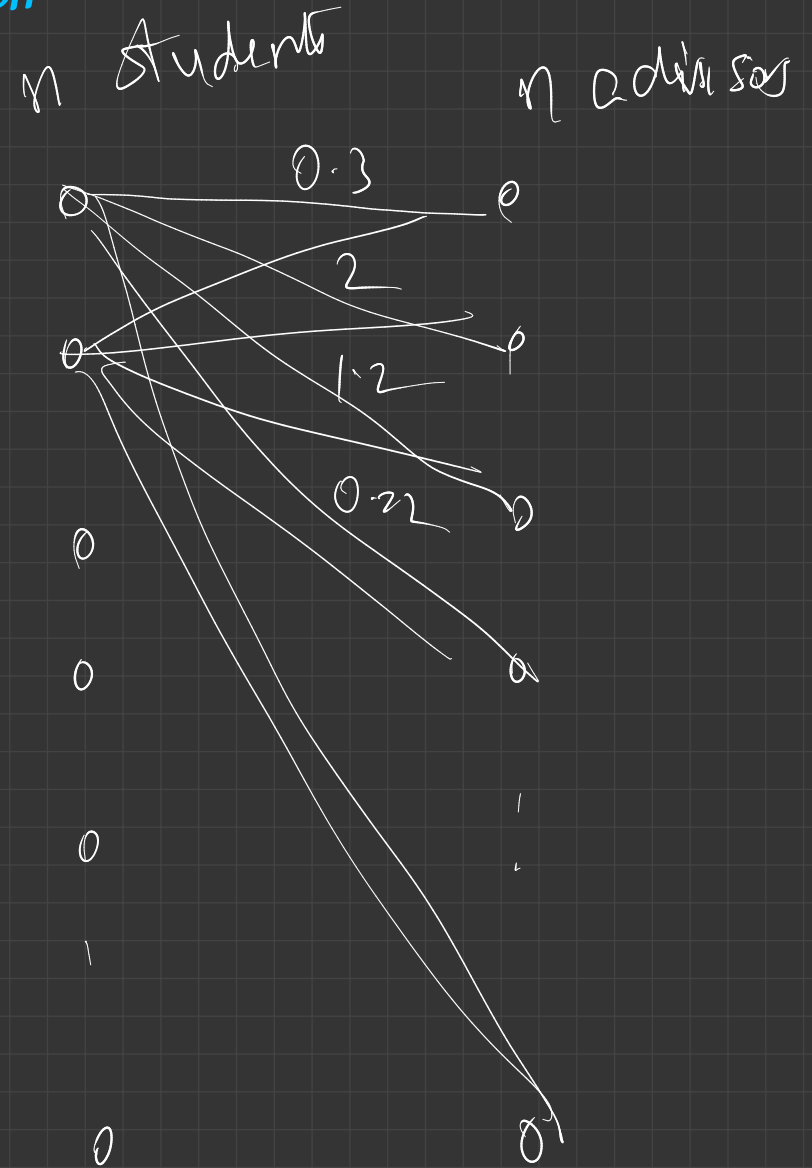
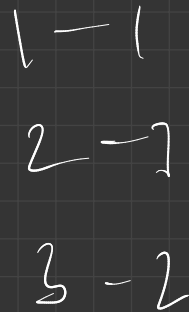
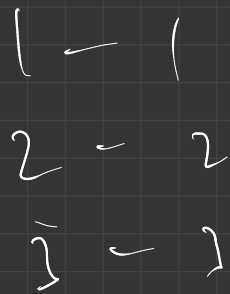
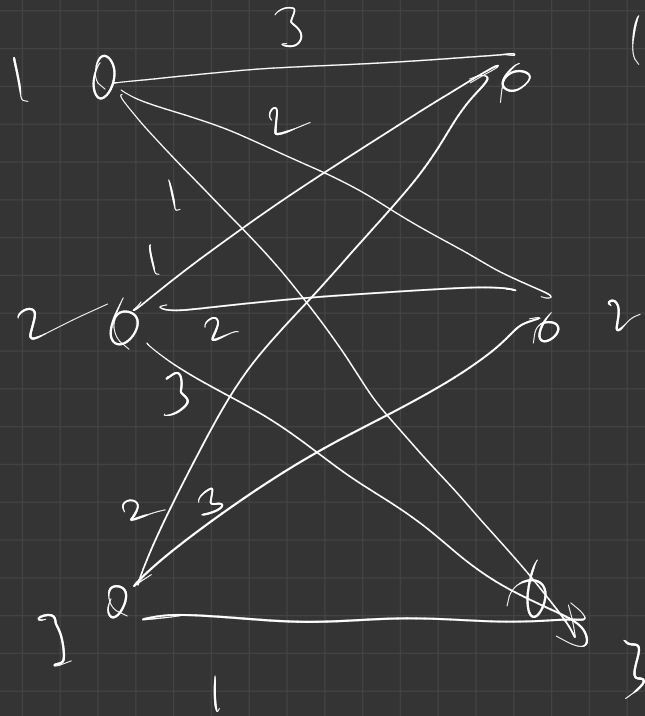
Polystope



convex hull of finitely
many pts

Two different ways of looking at a polytope

Maximum weight matching on a complete bipartite graph



Problem,

W is $n \times n$

w_{ij} is the score given by i about j

$$\begin{bmatrix} 3 & 2 & 1 \\ 1 & 2 & 3 \\ 2 & 3 & 1 \end{bmatrix}$$

Variables,

x is $n \times n$

$x_{ij} \in [0, 1]$

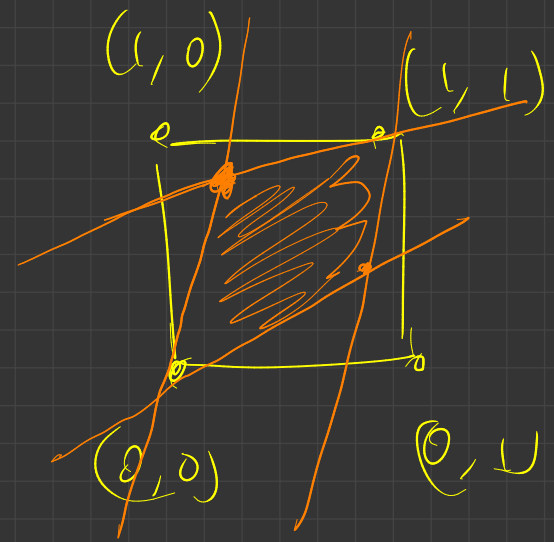
→ relax this to $x_{ij} \in [0, 1]$

$$\sum_j x_{ij} = 1 \quad \forall i$$

$$\sum_i x_{ij} = 1 \quad \forall j$$

constraints

Objective, $\max_x \sum_{i,j} w_{ij} x_{ij}$



Doubly stochastic matrices and the Birkhoff polytope

X is a doubly stochastic matrix if:

$$X \quad x_{ij} \in [0, 1]$$

$$\sum_j x_{ij} = 1 \quad \forall i, \quad \sum_i x_{ij} = 1 \quad \forall j$$

If $x_{ij} \in [0, 1]$ in addition, X is a permutation matrix.

* Any convex combination of permutation matrices is doubly stochastic

* Birkhoff-von Neumann theorem: Birkhoff polytope = conv(permutation matrices)

Norm balls and ellipsoids

$$\mathcal{B}_n(\underline{x}_c, n) = \{ \underline{x} \in \mathbb{R}^n, \|\underline{x} - \underline{x}_c\| \leq n \}$$

$$\underline{x}_1, \underline{x}_2 \in \mathcal{B}_n(\underline{x}_c, n)$$

$$\| \alpha \underline{x}_1 + (1-\alpha) \underline{x}_2 - \underline{x}_c \|$$

$$= \| \alpha(\underline{x}_1 - \underline{x}_c) + (1-\alpha)(\underline{x}_2 - \underline{x}_c) \|$$

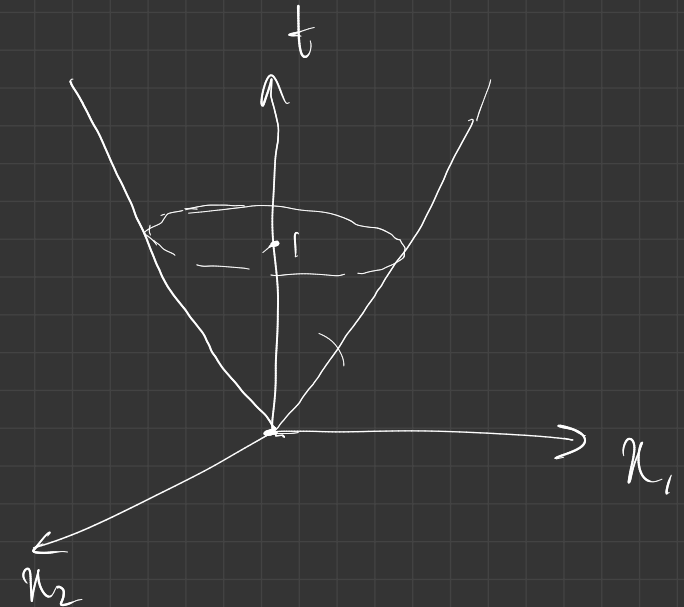
$$\leq \| \alpha(\underline{x}_1 - \underline{x}_c) \| + \| (1-\alpha)(\underline{x}_2 - \underline{x}_c) \|$$

$$= \alpha \| \underline{x}_1 - \underline{x}_c \| + (1-\alpha) \| \underline{x}_2 - \underline{x}_c \|$$

$$\leq \alpha n + (1-\alpha)n = n$$

Norm cones and the positive semidefinite cone

Norm cone : $\left\{ (\underline{x}, t) = \underline{x} \in \mathbb{R}^n, t \geq 0, \|\underline{x}\|_2 \leq t \right\}$



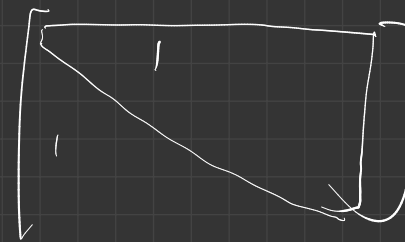
$$\alpha_1 \begin{bmatrix} \underline{x}_1 \\ t_1 \end{bmatrix} + \alpha_2 \begin{bmatrix} \underline{x}_2 \\ t_2 \end{bmatrix} \quad \alpha_1, \alpha_2 \geq 0$$

$$\begin{bmatrix} \alpha_1 \underline{x}_1 + \alpha_2 \underline{x}_2 \\ \alpha_1 t_1 + \alpha_2 t_2 \end{bmatrix}$$

$$\begin{aligned} \|\alpha_1 \underline{x}_1 + \alpha_2 \underline{x}_2\| &\leq \|\alpha_1 \underline{x}_1\| + \|\alpha_2 \underline{x}_2\| \\ &= \alpha_1 \|\underline{x}_1\| + \alpha_2 \|\underline{x}_2\| \\ &\leq \alpha_1 t_1 + \alpha_2 t_2 \end{aligned}$$

$$\Rightarrow \alpha_1 \begin{bmatrix} \underline{x}_1 \\ t_1 \end{bmatrix} + \alpha_2 \begin{bmatrix} \underline{x}_2 \\ t_2 \end{bmatrix} \in C$$

$$\mathcal{S}^n = \{ A \in \mathbb{R}^{n \times n} : A^T = A \}$$



$$\dim(\mathcal{S}^n) = n(n+1)/2$$

$$\mathcal{S}_+^n = \{ A \in \mathcal{S}^n, A \text{ is PSD} \}$$

Convex?

A_1, A_2 PSD

$$\underline{x}^T (\alpha_1 A_1 + \alpha_2 A_2) \underline{x} = \alpha_1 \underline{x}^T A_1 \underline{x} + \alpha_2 \underline{x}^T A_2 \underline{x} \geq 0$$

\mathcal{S}_+^n is a convex cone

Transformations of Convex Sets

1. Intersection of convex sets

A_1, A_2 are convex

$A_1 \cap A_2$ is also convex

$$\underline{x}_1, \underline{x}_2 \in A_1 \cap A_2$$

$$\underline{y} = \alpha \underline{x}_1 + (1-\alpha) \underline{x}_2$$

$$\underline{x}_1, \underline{x}_2 \in A_1 \Rightarrow \underline{y} \in A_1$$

$$\underline{x}_1, \underline{x}_2 \in A_2 \Rightarrow \underline{y} \in A_2$$

consider a family of sets $A_t : t \in \mathbb{R}$

$$\underline{x}_1, \underline{x}_2 \in A_t \quad \forall t$$

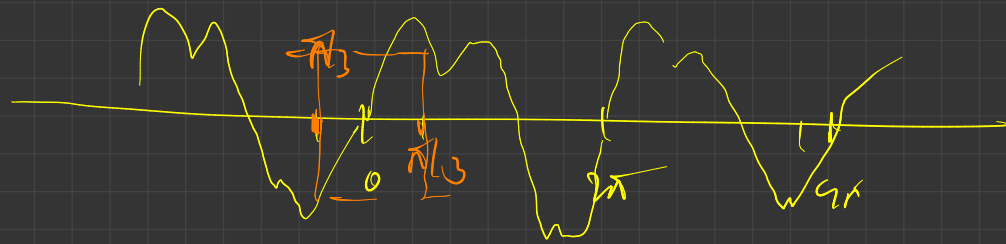
$$\underline{y} \in A_t \Rightarrow \underline{y} \in \bigcap_t A_t$$

Example :

$$\mathcal{A} = \{ \underline{x} \in \mathbb{R}^n \quad \& \}$$

$$\left. \left| \sum_{k=1}^n \underbrace{a_k}_{f_{\underline{x}}(t)} \cos(kt) \right| \leq 1 \right\}$$

$$\forall -\frac{\pi}{3} \leq t \leq \frac{\pi}{3}$$



$$\mathcal{A}_t = \left\{ \underline{x} \in \mathbb{R}^n : \left| \sum_{k=1}^n a_k \cos kt \right| \leq 1 \right\}$$

$$\langle \underline{x}, \underline{\cos kt} \rangle \leq 1$$

$$\langle \underline{x}, \underline{\cos kt} \rangle \geq -1$$

$$A = \bigcap_{t \in [-\pi/3, \pi/3]} A_t$$

$$A_t = \left\{ \underline{r} : \sum_{k=1}^n r_k e^{kt} \leq 1 \right\}$$

$$A_t = \left\{ (r_1, r_2) : \begin{aligned} r_1 \cos t + r_2 \cos 2t &\leq 1 \\ r_1 \cos t + r_2 \cos 2t &\geq -1 \end{aligned} \right\}$$

$$r_1 \leq \frac{1 - r_2 \cos 2t}{\cos t}$$

$$r_1 \geq \frac{-1 - r_2 \cos(2t)}{\cos t}$$

2. Minkowski sum of convex sets

A_1, A_2 convex

$$A_1 + A_2 = \{ \underline{x}_1 + \underline{x}_2 \mid \underline{x}_1 \in A_1, \underline{x}_2 \in A_2 \}$$

$$\begin{array}{l}
 \underline{x}_1, \underline{x}_2 \in A_1 + A_2 \\
 \parallel \\
 \underline{x}_1 + \underline{x}_2 \\
 \parallel \\
 \underline{x}_1 \in A_1 \quad \underline{x}_2 \in A_2 \\
 \parallel \\
 \alpha \underline{x}_1 + (1-\alpha) \underline{x}_2 \\
 \parallel \\
 \alpha \underline{x}_1 + \alpha \underline{x}_2 + (1-\alpha) \underline{x}_1 + (1-\alpha) \underline{x}_2 \\
 \parallel \\
 \in A_1 + A_2
 \end{array}$$

$\in A_1$ (under $\alpha \underline{x}_1$)
 $\in A_2$ (under $\alpha \underline{x}_2$)
 $\in A_1$ (under $(1-\alpha) \underline{x}_1$)
 $\in A_2$ (under $(1-\alpha) \underline{x}_2$)

3. Cartesian product of convex sets

A_1, A_2 are convex

$$A_1 \subseteq \mathbb{R}^n \quad A_2 \subseteq \mathbb{R}^m$$

$$A_1 \times A_2 = \{ (x, y) \mid x \in A_1, y \in A_2 \} \subseteq \mathbb{R}^{n+m}$$

4. Affine transform of a convex set

If A is convex, then (for $A \in \mathbb{R}^{m \times n}$, $\underline{b} \in \mathbb{R}^m$)

$A' = A A + \underline{b} = \{ A \underline{x} + \underline{b} : \underline{x} \in A \}$ is convex

Proper cone

A cone \mathcal{K} is proper if

- ① \mathcal{K} is closed
- ② \mathcal{K} is convex
- ③ \mathcal{K} is solid: has a nonempty interior
- ④ \mathcal{K} is pointed: $\underline{x} \in \mathcal{K}, \underline{x} \neq \underline{0} \Rightarrow -\underline{x} \notin \mathcal{K}$.

— Cone that is not closed:

$(\underline{x}_1, \underline{x}_2)$:

$$\underline{x}_1 > 0$$

$$\underline{x}_2 > 0$$

$$\text{or } \underline{x}_2 = \underline{x}_2 = 0 \}$$

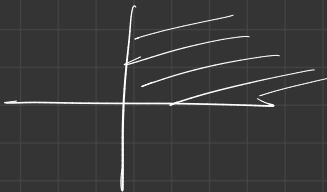


$$(\underline{x}_1, a\underline{x}_1)$$

- Convex not solid, ray in \mathbb{R}^2
 - Not solid
 - Convex
 - closed
 - pointed

- Convex not pointed: ① line passing through origin!
 - Not solid
 - Convex
 - closed
 - Not pointed

- ② \mathbb{R}^2 :
 - Solid
 - convex
 - closed
 - not pointed

Example: ① $\{ \underline{x} : x_i \geq 0 \ \forall i \}$ \rightarrow 

② The set of all PSD matrices $\subseteq \mathcal{S}^n$

(i) If A is PSD then αA is PSD for $\alpha \geq 0$

(ii) If $A_1, A_2 \in \mathcal{S}_+^n$ then $\alpha A_1 + (1-\alpha) A_2$ is PSD

$$\underline{x}^T (\alpha A_1 + (1-\alpha) A_2) \underline{x} \quad \text{for any } \underline{x} \in \mathbb{R}^n \\ \geq \alpha \underline{x}^T A_1 \underline{x} + (1-\alpha) \underline{x}^T A_2 \underline{x} \\ \geq 0 \quad \text{Hence convex} \\ 0 \leq \alpha \leq 1$$

(iii) $A \in \mathcal{S}_+^n$ $A \neq \underline{0}$ $-A$ is not PSD
 $-A \notin \mathcal{S}_+^n$

(iv) In \mathbb{R}^n $B_n(n) = \{ \underline{x} \in \mathbb{R}^n : \|\underline{x}\|_2^2 \leq n^2 \}$

$$\sum_{i=1}^n x_i^2$$

for S^n , (or even $\mathbb{R}^{m \times n}$), we use the Frobenius norm

$$\|A\| = \sqrt{\sum_{i,j} a_{ij}^2}$$

Consider $A = \begin{bmatrix} a_{11} & a_{12} \\ a_{12} & a_{22} \end{bmatrix}$

$$\|A\|^2 = a_{11}^2 + a_{22}^2 + 2a_{12}^2$$

↓
3 free variables

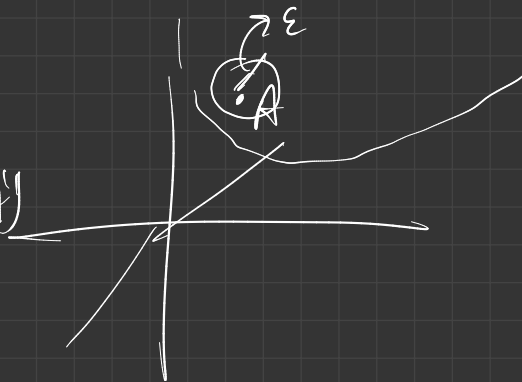
$$\begin{bmatrix} a_{11} \\ a_{12} \\ a_{22} \end{bmatrix}$$

Claim: $\text{int}(S_{++}^n) = S_{++}^n$

\equiv
 $\exists A \in S_{++}^n$ then \exists a ball of radius ε around A

\equiv
 Fix any $A \in S_{++}^n$. $\exists \varepsilon > 0$ $A \in S_{++}^n$
 st $A + A' \in S_{++}^n$ $\iff A' \in S^n$ $A' \in S^n$
 \cap
 S_{++}^n as long as $\|A'\| < \varepsilon$

$$\lambda_{\min}(A) + \lambda_{\min}(A') \leq \lambda_{\min}(A+A') \leq \lambda_{\max}(A+A') \leq \lambda_{\max}(A) + \lambda_{\max}(A')$$



As long as $\lambda_{\min}(A') > -\lambda_{\min}(A)$

$\lambda_{\min}(A+A') > 0 \implies A+A'$ is PD.

Generalized Inequalities

Given any proper cone K .

$$\underline{x} \preceq_K y \iff y - \underline{x} \in K$$

① K the nonnegative orthant (in \mathbb{R}^n)

$$K = \{ \underline{x} \in \mathbb{R}^n : x_i \geq 0 \forall i \}$$

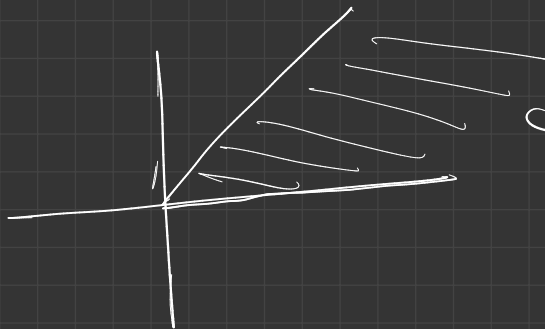
$$\underline{x} \preceq_K y \Rightarrow y - \underline{x} \in K$$

$$\Rightarrow y_i - x_i \geq 0 \quad \forall i$$

$$y_i \geq x_i \quad \forall i$$

②

$$K =$$

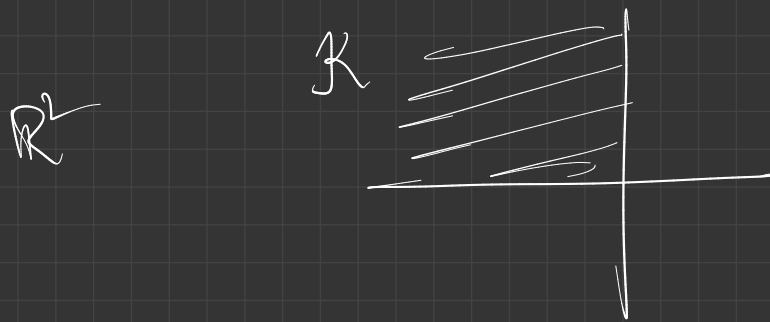


$$\rightarrow K = \{ (x, y) : x \geq 0, y \leq x \}$$

$$\begin{bmatrix} x_1 \\ y_1 \end{bmatrix} \preceq_K \begin{bmatrix} x_2 \\ y_2 \end{bmatrix} \iff \begin{bmatrix} x_2 - x_1 \\ y_2 - y_1 \end{bmatrix} \in K$$

$$\iff \begin{aligned} x_2 &\geq x_1 \\ y_2 - y_1 &\leq x_2 - x_1 \end{aligned}$$

②



$$K = \{(x, y) : \begin{aligned} x &\geq 0 \\ y &\geq 0 \end{aligned}\}$$

$$\begin{bmatrix} x_1 \\ y_1 \end{bmatrix} \preceq_K \begin{bmatrix} x_2 \\ y_2 \end{bmatrix} \iff \begin{bmatrix} x_2 - x_1 \\ y_2 - y_1 \end{bmatrix} \in K$$

$$\iff \begin{aligned} x_2 - x_1 &\geq 0 & \& \quad y_2 - y_1 &\geq 0 \end{aligned}$$

$$\iff \begin{aligned} x_2 &\geq x_1 & \& \quad y_2 &\geq y_1 \end{aligned}$$

$$\textcircled{1} \quad \mathcal{K} = \mathbb{S}_+^n$$

$$A \preceq_{\mathcal{K}} B \iff B - A \text{ is PSD}$$

Properties:

$$\textcircled{1} \quad x \preceq_{\mathcal{K}} y \quad \text{then} \quad x+z \preceq_{\mathcal{K}} y+z \quad \forall z \in \mathbb{R}^n$$

$$\textcircled{2} \quad x \preceq_{\mathcal{K}} y \quad \wedge \quad y \preceq_{\mathcal{K}} z \quad \text{then} \quad x \preceq_{\mathcal{K}} z$$

$$\Downarrow$$

$$y - x \in \mathcal{K}$$

$$\Downarrow$$

$$z - y \in \mathcal{K}$$

Since \mathcal{K} is a convex cone, $(y-x) + (z-y) \in \mathcal{K}$

$$\Rightarrow z - x \in \mathcal{K}$$

$$\Rightarrow x \preceq_{\mathcal{K}} z$$

$$\textcircled{3} \quad \underline{a} \leq_K y \quad \Rightarrow \quad \alpha \underline{a} \leq_K \alpha y \quad \forall \alpha \geq 0$$

$$y - \underline{a} \in K \quad \Rightarrow \quad \alpha(y - \underline{a}) \in K$$

$$\textcircled{4} \quad \underline{a} \leq_K \underline{a} \quad \text{sim} \quad \underline{0} \in K$$

$$\textcircled{5} \quad \underline{a} \leq_K y \quad \wedge \quad y \leq_K \underline{a} \quad \Rightarrow \quad \underline{a} = y$$

$$y - \underline{a} \in K \quad \wedge \quad (\underline{a} - y) = -(y - \underline{a}) \in K$$

$$\Rightarrow y = \underline{a}$$

$$\textcircled{6} \quad \text{If} \quad \underline{a}_n \leq_K y_n \quad \text{for} \quad n=1, 2, 3, 4, \dots$$

$$\lim_{n \rightarrow \infty} \underline{a}_n \leq_K \lim_{n \rightarrow \infty} y_n$$

If

$$x_n \leq y_n \quad \forall n$$

then

is

$$\lim_{n \rightarrow \infty} x_n \leq \lim_{n \rightarrow \infty} y_n \quad ?$$

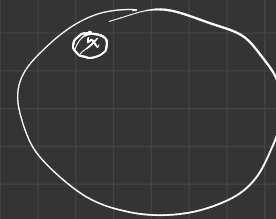
$$x_n = \frac{1}{n}$$

$$y_n = \frac{2}{n}$$

Want to show S_{++}^n is open in S^n

show for any $A \in S_{++}^n$, a ball of radius $\varepsilon \ll 1$
centered at A is contained in S_{++}^n

$A+B \in S_{++}^n$ for any B
 $\|B\| < \varepsilon$



Example: If \underline{x} is a random vector

Correlation matrix: $\mathbb{E}[\underline{x} \underline{x}^T] = C_x$

If x, y rv's

Correlation matrix $= \begin{bmatrix} \mathbb{E}x^2 & \mathbb{E}xy \\ \mathbb{E}xy & \mathbb{E}y^2 \end{bmatrix}$

$$\begin{aligned} \underline{u}^T C_x \underline{u} &= \underline{u}^T \mathbb{E}[\underline{x} \underline{x}^T] \underline{u} \\ &= \mathbb{E}[\underline{u}^T \underline{x} \underline{x}^T \underline{u}] \\ &= \mathbb{E}[(\underline{u}^T \underline{x})^2] \geq 0 \end{aligned}$$

X, Y, Z cols

$$C_X = \begin{bmatrix} 1 & \rho_{xy} & \rho_{xz} \\ \rho_{xy} & 1 & \rho_{yz} \\ \rho_{xz} & \rho_{yz} & 1 \end{bmatrix}$$

Given:

$$-0.2 \leq \rho_{xy} \leq 0.3$$

$$\rho_{xz} \geq 0$$

What is max ρ_{yz} ?

Problem:

$$\max \rho_{yz} \quad \text{st} \quad -0.2 \leq \rho_{xy} \leq 0.3$$

$$\rho_{xz} \geq 0$$

$$\& C_X \text{ is PSD.}$$

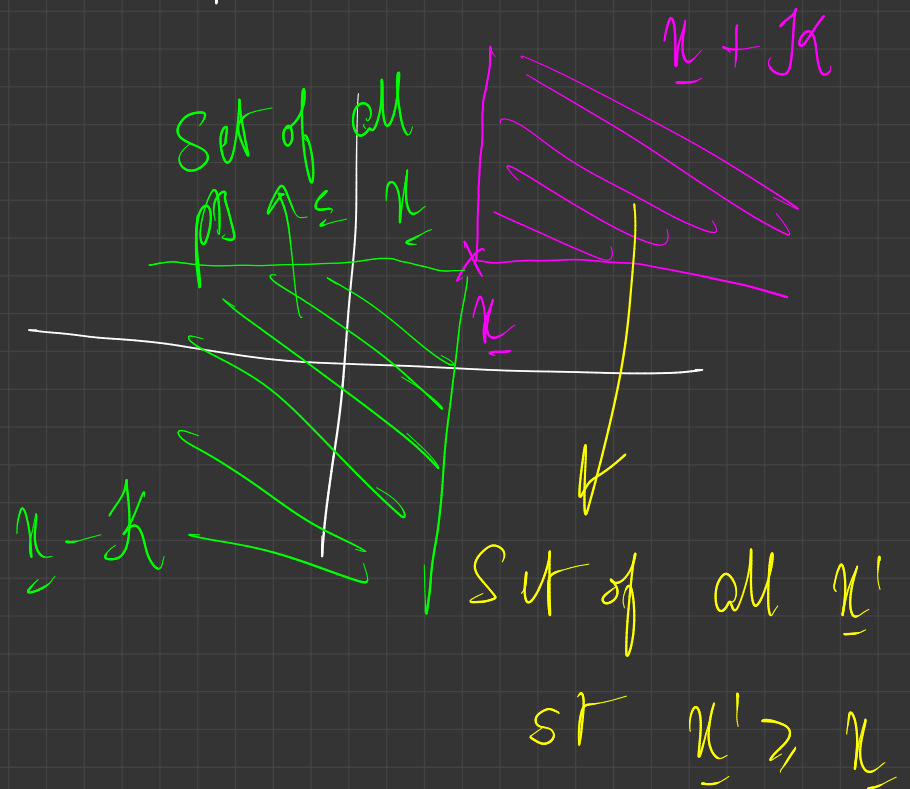
Generalized inequalities only form a partial order

$$\begin{bmatrix} 1 \\ 2 \end{bmatrix} \leq \begin{bmatrix} 3 \\ 4 \end{bmatrix}$$

$$\begin{bmatrix} 1 \\ 2 \end{bmatrix} \not\leq \begin{bmatrix} 2 \\ 1 \end{bmatrix}$$

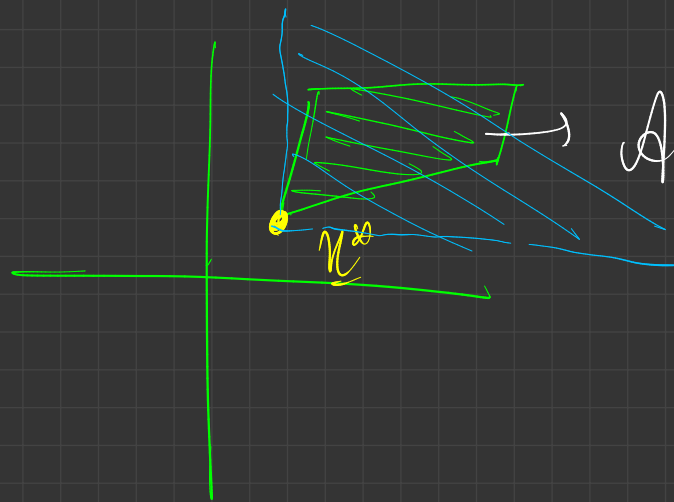
Not all pairs of elements are comparable.

Componentwise inequality



Componentwise inequality:

Does A have
a minimum?



$$\underline{x}^* = \min A \quad \text{iff}$$

$$\underline{x}^* \leq \underline{x} \quad \forall \underline{x} \in A$$

$$\underline{x}^* \in A$$

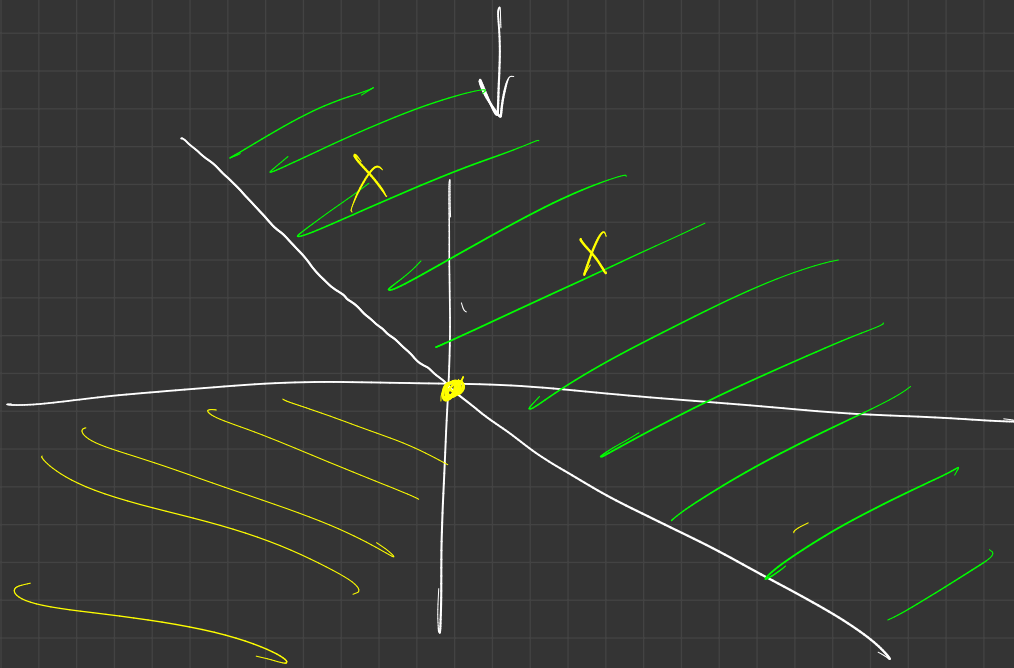
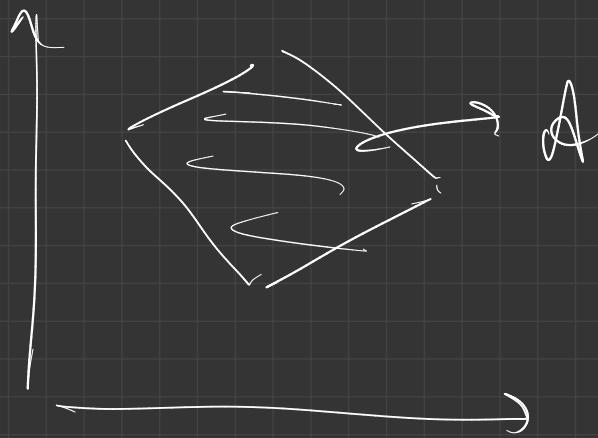
Observation:

$$\underline{x}^* = \min A$$

$$\text{iff } (\underline{x}^* + \mathcal{K}) \supseteq A$$

$$\hookrightarrow \underline{x}^* \in A$$

② $f(x_1, x_2) = x_1 + x_2 \geq 0$



Minimal points
of d

We say that $\underline{x} \in d$ is a minimal pt
if $y \leq \underline{x}$ & $y \in d \Rightarrow y = \underline{x}$

Similarly $\underline{n} \in A$ is a maximal pt of A if

$$y \geq \underline{n} \wedge y \in A \Rightarrow y = \underline{n}$$

Separating hyperplanes

Given A, B . We say that $\{ \underline{a}^T \underline{x} = b \}$

is a separating hyperplane for A, B if

$$\underline{a}^T \underline{x} \geq b \quad \forall \underline{x} \in A$$

$$\underline{a}^T \underline{x} \leq b \quad \forall \underline{x} \in B$$

The hyperplane strictly separates A, B if

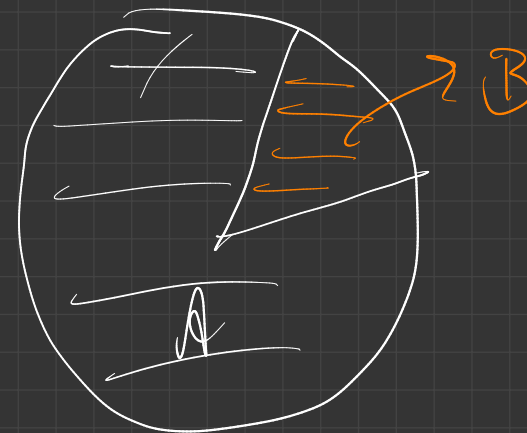
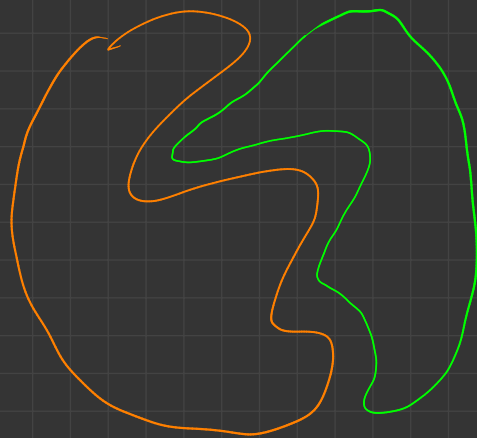
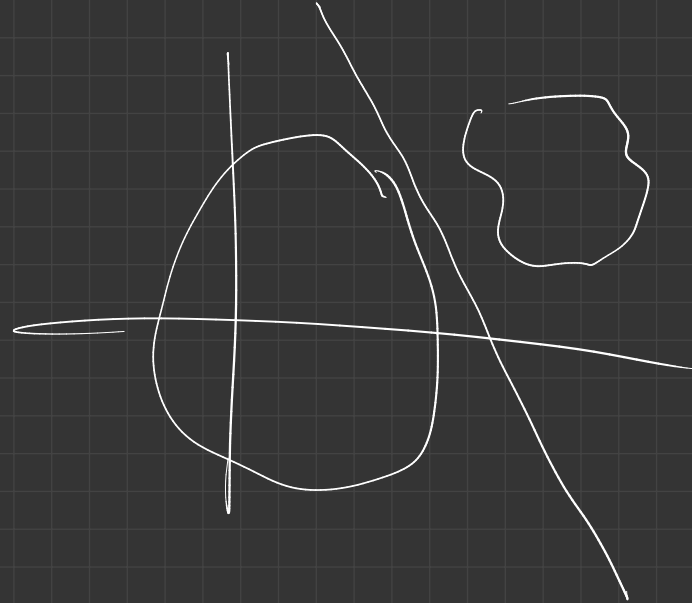
$$\underline{a}^T \underline{x} > b \quad \forall \underline{x} \in A$$

$$\underline{a}^T \underline{x} < b \quad \forall \underline{x} \in B.$$

Theorem

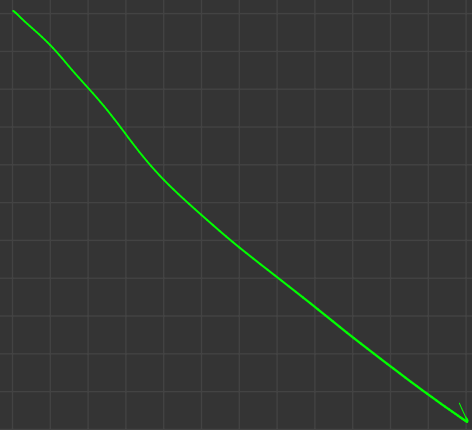
If A, B convex
& $A \cap B = \emptyset$, then

there is a hyperplane
that separates A & B .



$$A = \{ \underline{x} : \underline{a}^T \underline{x} \geq b \}$$

$$B = \{ \underline{x} : \underline{a}^T \underline{x} \leq b \}$$



* Convex not true in general

* Convex true with additional constraints

- Strict separation, or

- One of A, B is open

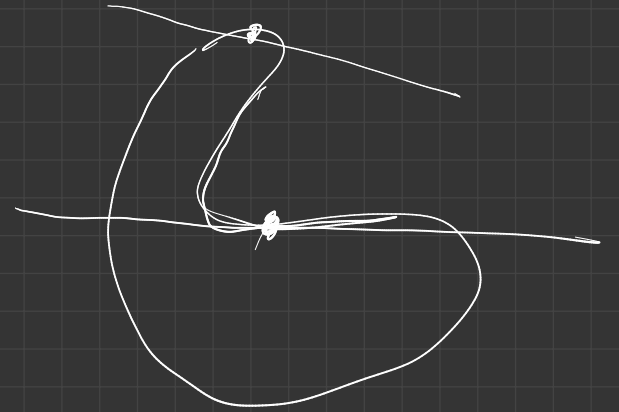
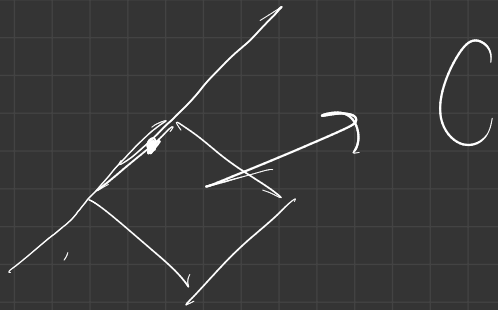
Supporting hyperplane

$$\underline{n}_0 \in \text{bd}(A) \xrightarrow{\text{boundary}} \partial(A) \setminus \text{int}(A)$$

We say that $\{ \underline{n} : \underline{a}^\top \underline{n} = \underline{a}^\top \underline{n}_0 \}$ is a supporting hyperplane for A if all of A lies on one side

$$\underline{a}^\top \underline{n} \geq \underline{a}^\top \underline{n}_0 \quad \forall \underline{n} \in A$$

Claim: Every pt on the boundary of a convex set has a supporting hyperplane



Take $\{x_0\} = A$

$B = \text{int}(C)$

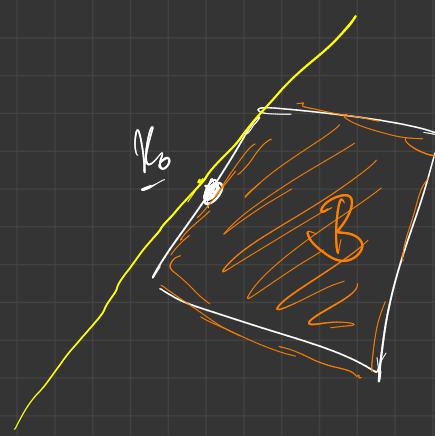
$A \cap B = \emptyset$

A, B convex

$\exists \underline{a} \in \mathbb{R}^n, b \in \mathbb{R}$ st $\underline{a}^T \underline{x} \geq b \quad \forall \underline{x} \in B$

$\underline{a}^T \underline{x} \leq b \quad \forall \underline{x} \in A$

\underline{x}_0 is on the
separating
hyperplane

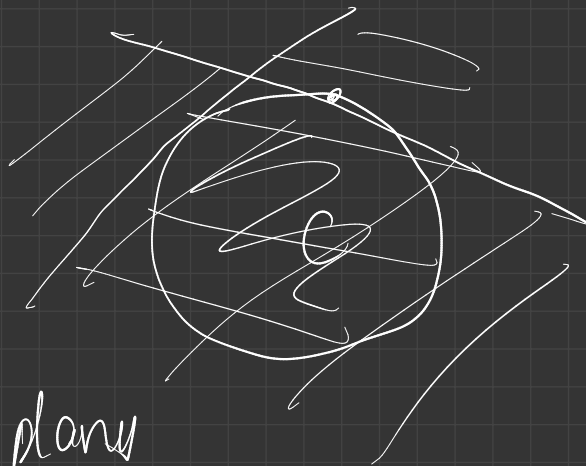


$$\Rightarrow \underline{a}^T \underline{x} \geq b$$

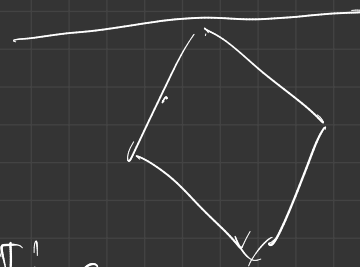
$$\forall \underline{x} \in C$$

\therefore this is a supporting hyperplane.

Every closed convex set is
the intersection of half spaces
defined by the supporting hyperplanes

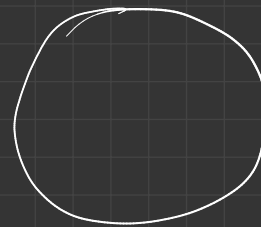
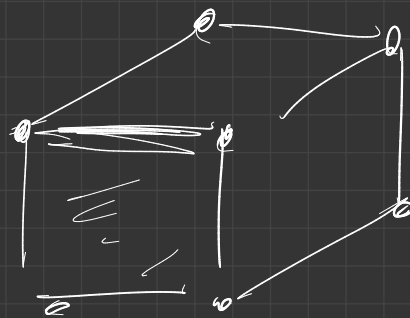


Extreme point: $\underline{x} \in A$ is an
extreme pt of supporting
hyperplane has only one pt from A
(\underline{x} itself)



Property :

A closed convex set
is the convex hull
of its extreme points



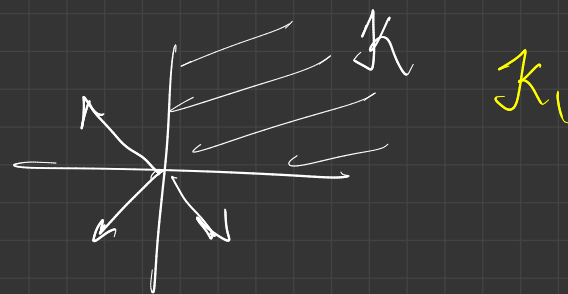
Dual cone

If K is a cone. The dual cone

$$K^{\circ} = \{ \underline{x} \in \mathbb{R}^n : \underline{x}^T \underline{y} \geq 0 \ \forall \underline{y} \in K \}$$

① Nonnegative orthant

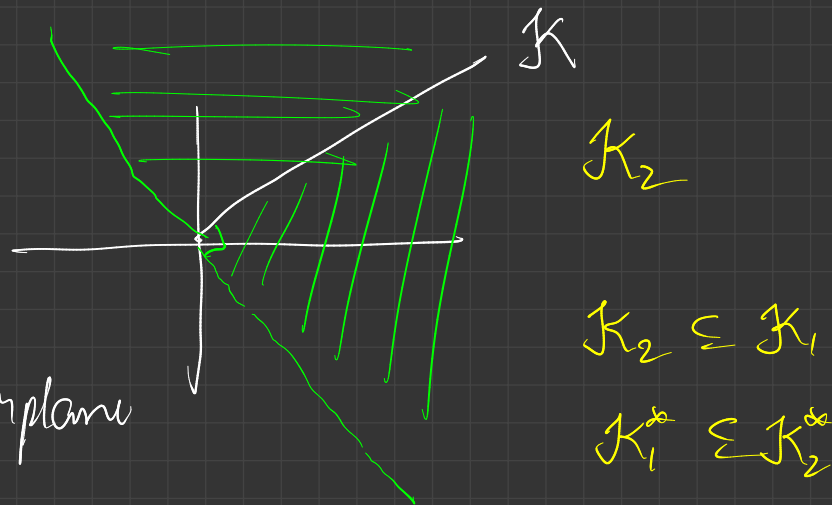
$$K^{\circ} = K$$



② $K = \{ \alpha \underline{x} : \alpha \geq 0 \}$

K° = half space defined by hyperplane normal to \underline{x}

$$= \{ \underline{y} : \underline{x}^T \underline{y} \geq 0 \}$$



③ $\mathcal{K} =$ any k -dim subspace of \mathbb{R}^n

$$\underline{x} \in \mathcal{K} \quad \underline{y} \in \mathcal{K}^*$$

Suppose $\underline{x}^T \underline{y} > 0$ $(-\underline{x})^T \underline{y} < 0$

$$\Rightarrow \underline{x}^T \underline{y} = 0 \quad \forall \underline{y} \in \mathcal{K}^*$$

$\mathcal{K}^* = \mathcal{K}^\perp$, the dual space of \mathcal{K}

④ $\mathcal{K} = \mathbb{S}_+^n \quad \langle A, B \rangle = \sum_{i,j} a_{ij} b_{ij} = \text{tr}(A^T B)$

Claim: $\mathcal{K}^* = \mathcal{K}$

Take $A \in \mathcal{K}^*$

want to show
 $\underline{x}^T A \underline{x} \geq 0$

$$\underline{y}^T (\underline{x} \underline{x}^T) \underline{y} = (\underline{y}^T \underline{x}) (\underline{x}^T \underline{y}) = (\underline{x}^T \underline{y})^2 \geq 0$$

$$\Rightarrow \underline{x} \underline{x}^T \text{ is PSD}$$

$$A \in \mathcal{K}^* \Rightarrow \underline{x}^T A \underline{x} = \underbrace{\text{tr}(\underline{x} \underline{x}^T A)}_{\geq 0} \geq 0$$

since $A \in \mathcal{K}^*$ & $\underline{x} \underline{x}^T \in \mathcal{K}$

Consider any PSD matrix A, B

$$A = \sum_{i=1}^n \lambda_i q_i q_i^T$$

$$\begin{aligned} \text{tr}(A^T B) &= \text{tr}\left(\sum_{i=1}^n \lambda_i q_i q_i^T B\right) = \sum_{i=1}^n \text{tr}(\lambda_i q_i q_i^T B) \\ &= \sum_{i=1}^n \lambda_i \text{tr}(q_i q_i^T B) \\ &= \sum_{i=1}^n \lambda_i \underbrace{(q_i^T B q_i)}_{\geq 0} \end{aligned}$$

$$\Rightarrow A \in \mathcal{K}^*$$

$$\Rightarrow \mathcal{K}^* = \mathcal{K}$$

Properties of dual cones

① K^* is closed & convex

$$r_1, r_2 \in K^*,$$

$$r_1^T y \geq 0 \quad \forall y \in K$$

$$r_2^T y \geq 0 \quad \forall y \in K$$

$$(\alpha r_1 + (1-\alpha)r_2)^T y \geq 0 \quad \forall y \in K$$

$$\Rightarrow \alpha r_1 + (1-\alpha)r_2 \in K^*$$

② $K_1 \subseteq K_2 \Rightarrow K_1^* \supseteq K_2^*$

Take any $r \in K_2^*$

$$r^T y \geq 0 \quad \forall y \in K_2$$

$$\Rightarrow r^T z \geq 0 \quad \forall z \in K_1$$

$$\Rightarrow r \in K_1^*$$

③ If K has nonempty interior, then K^* is pointed.

Proof: Suppose K^* is not pointed,

$$\exists \text{ nonzero } \underline{u} \in K^* \text{ or } -\underline{u} \in K^*$$

$$\text{Take any } y \in K \quad y^T \underline{u} = 0$$

$$\dim(K) \leq n-1$$

$\Rightarrow K$ does not have a nonempty interior

closure of

④ If K is pointed, K^* has nonempty interior

Proof: Suppose K^* does not have a nonempty interior

$$\Leftrightarrow \dim(K^*) \leq n-1$$

$$\Rightarrow \exists y \in \mathbb{R}^n \text{ or } y^T \underline{u} = 0 \quad \forall \underline{u} \in K^*$$

$$\Rightarrow y \in K \text{ \& } K \text{ is not pointed,}$$

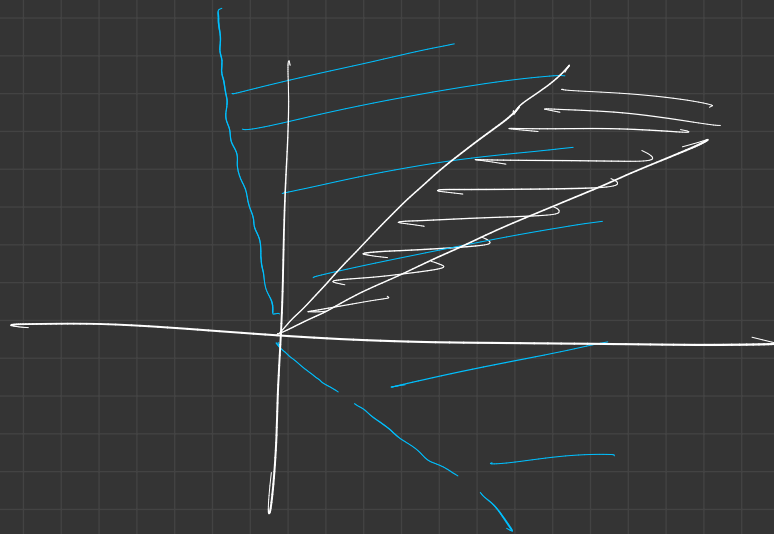
K^* is not pointed

But K^{**} is not pointed $\Rightarrow d(K)$ is not pointed

(If closure of K is pointed, then the convex hull of the closure of K cannot be pointed)

However, $K = \{0\} \cup \{(x_1, x_2) : x_1 \in \mathbb{R}, x_2 > 0\}$ is a pointed cone, but closure is not pointed. - The closure is $\{(x_1, x_2) : x_1 \in \mathbb{R}, x_2 \geq 0\}$ & $(-1, 0)$ & $(1, 0)$ lie in this set.

The dual of K is $\{(0, x_2) : x_2 \geq 0\}$. This has empty interior.



Property: If \mathcal{K} is a proper cone, then \mathcal{K}^* is a proper cone

$$\begin{aligned} x \preceq_{\mathcal{K}} y &\iff y - x \in \mathcal{K} \\ y \succeq_{\mathcal{K}} 0 &\implies y \in \mathcal{K} \end{aligned}$$

PROP: $x \preceq_{\mathcal{K}} y \iff \left(\lambda \in \mathcal{K}^* \implies \lambda \succeq_{\mathcal{K}^*} 0 \right)$

$$\implies \lambda^T x \leq \lambda^T y$$

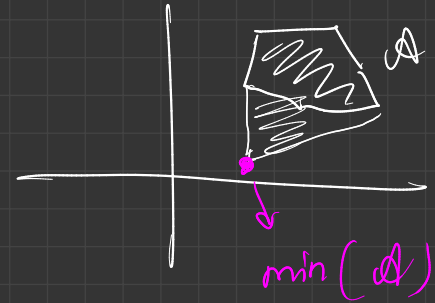
$$y - x \in \mathcal{K} \implies \lambda^T (y - x) \geq 0$$

$$\implies \lambda^T y \geq \lambda^T x$$

Consider $A \subseteq \mathbb{R}^n$ \mathcal{K} a proper cone

$$\underline{x}^* = \min(A)$$

Claim: $\underline{x}^* = \operatorname{argmin}_{x \in A} \lambda^\top x$ for every $\lambda \succ_{\mathcal{K}} 0$



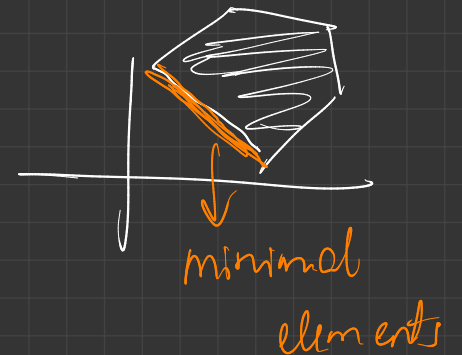
Proof: $\underline{x}^* \preceq_{\mathcal{K}} x \quad \forall x \in A$

$$\Rightarrow \lambda^\top \underline{x}^* \leq \lambda^\top x \quad \forall x \in A$$

Claim: If $\lambda \succ_{\mathcal{K}} 0$ \checkmark

$\underline{x}^* = \operatorname{argmin}_{x \in A} \lambda^\top x$, then \underline{x}^*

is minimal.



Proof: If not, $\exists x \in A \quad x^* \succ_{\mathcal{K}} x \Rightarrow \lambda^\top x^* > \lambda^\top x$

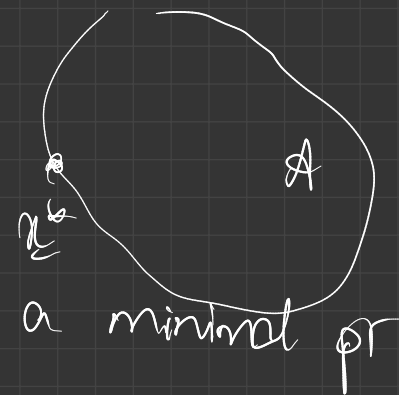
Consider the componentwise inequality

$x \in A$ is a minimum vector

$$\begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix}$$

$$\begin{bmatrix} x_1^* \\ x_2^* \\ x_3^* \end{bmatrix}$$

$$\begin{bmatrix} x_1^* \\ x_2^* \\ x_3^* \end{bmatrix} \in A$$



A minimal point is called a Pareto optimal point

$P_1 \backslash P_2$	C	NC
C	(2, 2)	(0, 4)
NC	(4, 0)	(1, 1)

dominant strategy but not Pareto optimal

all Pareto optimal

Problem: $\{ \underline{x} : \underline{a}_1^T \underline{x} = b_1 \} = \mathcal{H}_1$

$$\{ \underline{x} : \underline{a}_2^T \underline{x} = b_2 \} = \mathcal{H}_2$$

When are these parallel?

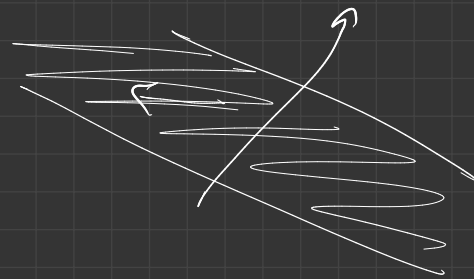
$$\underline{x}_1, \underline{x}_2 \in \mathcal{H}_1$$

$$\underline{a}_1^T (\underline{x}_2 - \underline{x}_1) = 0$$

→ vector that is "along" \mathcal{H}_1

\underline{a}_1 is normal to \mathcal{H}_1

\mathcal{H}_1 & \mathcal{H}_2 are parallel if $\underline{a}_1 = \alpha \underline{a}_2$ for some nonzero α

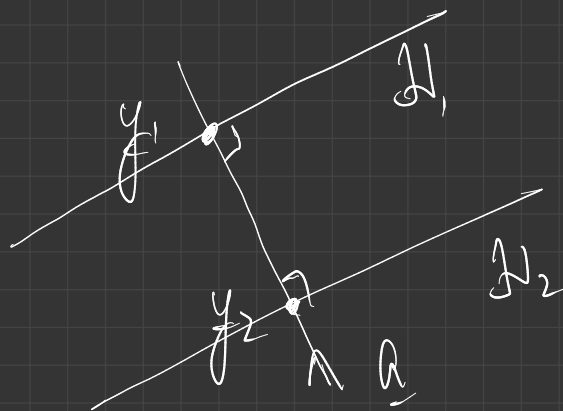


Suppose $\mathcal{H}_1 \parallel \mathcal{H}_2$. What is the distance b/w \mathcal{H}_1 & \mathcal{H}_2 ?

$$y_1 \in \text{ray}(\underline{a})$$

$$y_1 = \alpha_1 \underline{a}$$

$$\begin{aligned} \& y_1 \in \mathcal{H}_1 &\Rightarrow & \underline{a}^\top y_1 = b_1 \\ & &\Rightarrow & \alpha_1 = \frac{b_1}{\|\underline{a}\|^2} \end{aligned}$$



$$y_1 = \frac{b_1}{\|\underline{a}\|^2} \underline{a}$$

$$y_2 = \frac{b_2}{\|\underline{a}\|^2} \underline{a}$$

$$\|y_2 - y_1\| = \frac{|b_2 - b_1|}{\|\underline{a}\|^2} \|\underline{a}\| = \frac{|b_2 - b_1|}{\|\underline{a}\|}$$

Given $C \subseteq \mathbb{R}^n$, finite

$\underline{x} \in C \rightarrow$ transmitted

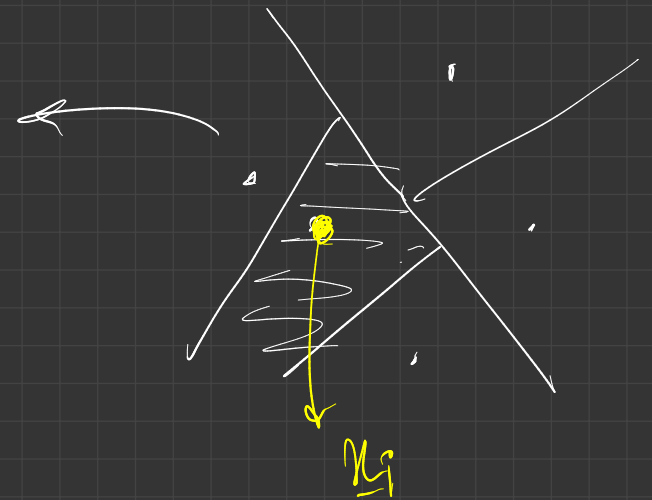
$\underline{y} = \underline{x} + \underline{e} \rightarrow$ received

Decision rule: Given $\underline{y} \in \mathbb{R}^n$,
choose $\hat{\underline{x}} = \operatorname{argmin}_{\underline{x} \in C} \|\underline{y} - \underline{x}\|$

$V_i = \{ \underline{y} \in \mathbb{R}^n :$

$\|\underline{y} - \underline{x}_i\| \leq \|\underline{y} - \underline{x}_j\|$ region
 $\forall j \neq i \}$

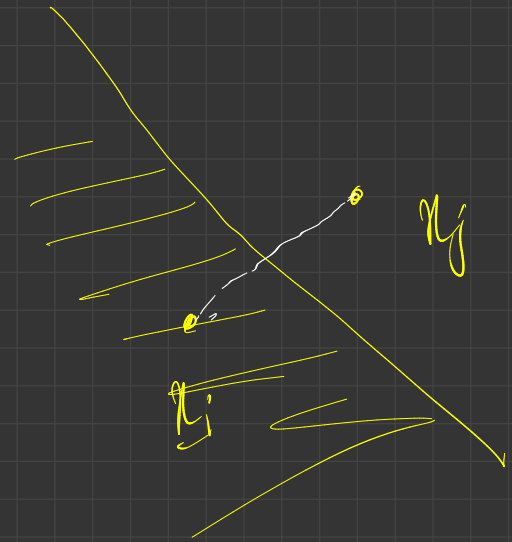
Voronoi



Voronoi region is a polyhedron

Let \mathcal{H} be the set of all pts in \mathbb{R}^n
which are closer to \underline{x}_i than \underline{x}_j .

Show that this is a half space.



$$\text{If } \underline{x} \in \mathcal{H}, \quad \|\underline{x} - \underline{x}_i\|^2 \leq \|\underline{x} - \underline{x}_j\|^2$$

$$(\underline{x} - \underline{x}_i)^T (\underline{x} - \underline{x}_i) \leq (\underline{x} - \underline{x}_j)^T (\underline{x} - \underline{x}_j)$$

$$\cancel{\underline{x}^T \underline{x}} - 2\underline{x}_i^T \underline{x} + \|\underline{x}_i\|^2 \leq \cancel{\underline{x}^T \underline{x}} - 2\underline{x}_j^T \underline{x} + \|\underline{x}_j\|^2$$

$$(\underline{x}_i - \underline{x}_j)^T \underline{x} \geq \frac{\|\underline{x}_i\|^2 - \|\underline{x}_j\|^2}{2}$$

Let H_{ij} be the halfspace of points in \mathbb{R}^n that are closer to x_i than x_j

$$H_{ij} = \{ \underline{x} \in \mathbb{R}^n : \| \underline{x} - x_i \| \leq \| \underline{x} - x_j \| \}$$

$$V_i = \{ \underline{x} \in \mathbb{R}^n : \| \underline{x} - x_i \| \leq \| \underline{x} - x_j \| \forall j \neq i \}$$

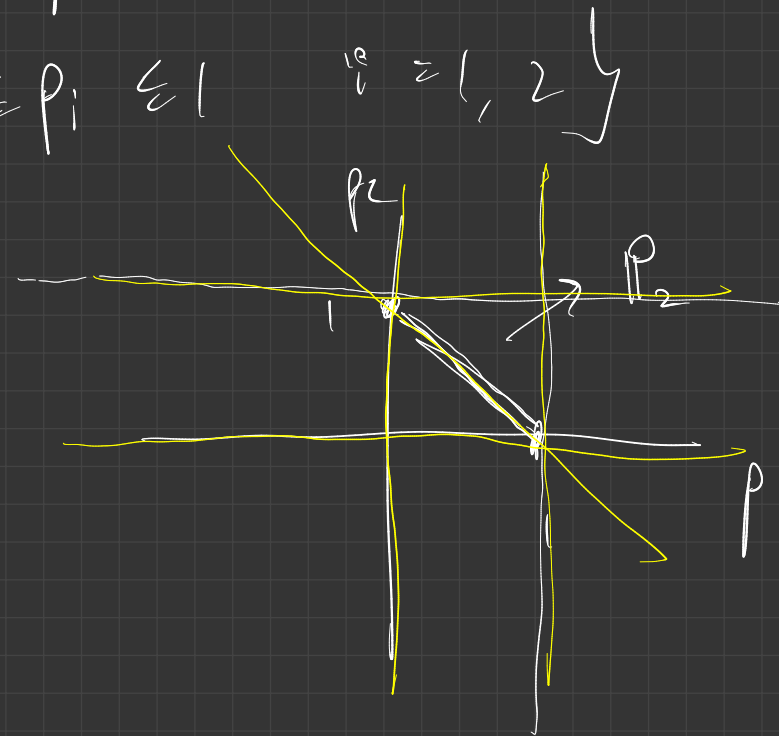
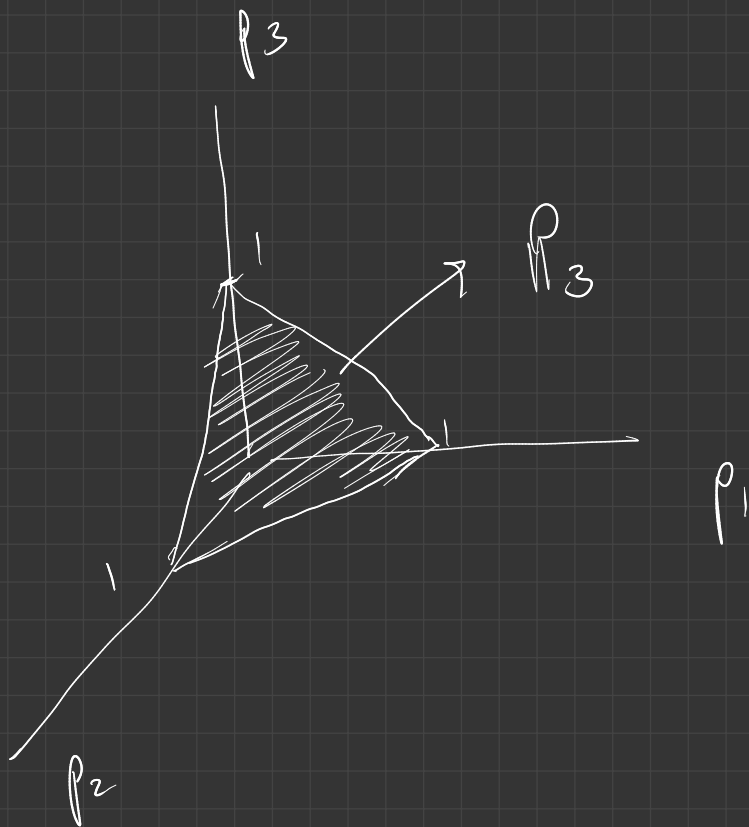
$$= \bigcap_{j \neq i} H_{ij}$$

$\therefore V_i$ is a polyhedron

③ PMF on $[n] = \{1, 2, 3, \dots, n\}$

$\mathcal{P}_n =$ set of all PMFs on $[n]$

$$\mathcal{P}_2 = \left\{ f \in \mathbb{R}^2 : \begin{array}{l} p_1 + p_2 = 1 \\ 0 \leq p_i \leq 1 \quad i = 1, 2 \end{array} \right\}$$



$$H_i = \left\{ p \in \mathbb{R}^n : \begin{array}{l} p_i \leq 1 \\ e_i^T p \leq 1 \end{array} \right\} \quad n \rightarrow \text{Halbspan}$$

$$\left\{ p \in \mathbb{R}^n : p_i \geq 0 \right\} \quad n \rightarrow \text{Halbspan}$$

$$\left\{ p \in \mathbb{R}^n : \sum_{i=1}^n p_i = 1 \right\} \quad 2$$

P_n is a polytope.

$$\mathcal{A} \approx \left\{ p \in \mathbb{P}_n : \underbrace{\mathbb{E} X^2 \leq \alpha, \quad X \sim p}_{\parallel} \right\}$$

$$\sum_{i=1}^n i^2 p_i \leq \alpha$$

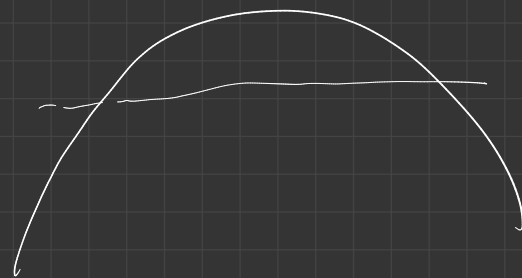
$$[1 \quad 2^2 \quad 3^2 \quad \dots \quad n^2] p \leq \alpha$$

$\Rightarrow \mathcal{A}$ is a polytope.

$$A = \left\{ f \in \mathbb{P}_n \mid H(f) \geq \alpha \right\}$$

$$H(f) = \sum_{i=1}^n p_i \log \frac{1}{p_i}$$

$\left\{ \underline{x} \mid f(\underline{x}) \geq \alpha \right\}$ is a convex set if f is concave



$\left\{ \underline{x} \mid f(\underline{x}) \leq \alpha \right\}$ is a convex set if f is convex

① Consider the set of all copositive matrices

$$C = \{ A \in S^n : \underline{x}^T A \underline{x} \geq 0 \text{ for all } \underline{x} \geq \underline{0} \}$$

Q1: Is this a cone?

$$\text{If } A \in C, \text{ then } \alpha A \in C \text{ for } \alpha \geq 0$$

Q2: Is this convex?

Q3: Is C closed?

The intersection of any number of closed sets is closed.

Fix an $\underline{x} \geq 0$. $\mathcal{H}_{\underline{x}} = \{ A : \underline{x}^T A \underline{x} \geq 0 \}$

↓

linear in A

$$\text{tr}(\underline{x} \underline{x}^T A) \geq 0$$

$\mathcal{H}_{\underline{x}}$ is a halfspace for any \underline{x}

$C = \bigcap_{\underline{x} \geq 0} \mathcal{H}_{\underline{x}} \rightarrow$ intersection of infinitely many halfspaces

Q4

Q4: Does C have nonempty interior?

$C \supseteq \mathcal{S}_+^n$ & \mathcal{S}_+^n has a nonempty interior

Q5: Is C pointed?

Yes

$A \in C \Rightarrow -A \notin C$ (if $A \neq 0$)

C is a proper cone.

HW: Find the dual cone of C .