

Advanced Matrix Algebra and Applications

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Books :

- 1) Fundamentals of matrix computations
— David Watkins
- 2) Matrix Algebra ...
— S. Boyd

Lecture 1 :

Inner product. (Dot product)

$$\mathbb{R}^n / \mathbb{R}^2 / \mathbb{R}^3$$

$$x, y \in \mathbb{R}^n$$

$$x = \begin{pmatrix} x_1 \\ \vdots \\ x_n \end{pmatrix}, \quad y = \begin{pmatrix} y_1 \\ \vdots \\ y_n \end{pmatrix}$$

$$\langle x, y \rangle = x^T y = \sum_{i=1}^n x_i y_i$$

$$\langle x, x \rangle = x^T x = \sum_{i=1}^n x_i^2 = \|x\|_2^2$$

$$\cos \theta = \frac{\langle x, y \rangle}{\|x\| \|y\|}$$

C.S. inequality $\Rightarrow -1 \leq \cos \theta \leq 1$

$x, y \in \mathbb{R}^n$ are such that

$$\langle x, y \rangle = x^T y = 0$$

LU
Cholesky
QR
SVD x

$Ax = b$ System of
linear equations.

$$A \in \mathbb{R}^{m \times n}$$

$$b \in \mathbb{R}^m$$

Q1: Does there exist $x \in \mathbb{R}^n$ s.t.
 $Ax = b$??

$$A = \begin{bmatrix} | & | & & | \\ A_1 & A_2 & \dots & A_n \\ | & | & & | \end{bmatrix}; \quad x = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

$$Ax = \begin{bmatrix} A_1 & \dots & A_n \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = x_1 A_1 + x_2 A_2 + \dots + x_n A_n$$

$$\text{Col span}(A) = \left\{ \alpha_1 A_1 + \alpha_2 A_2 + \dots + \alpha_n A_n \mid \right. \\ \left. \alpha_1, \dots, \alpha_n \in \mathbb{R} \text{ and } A_1, \dots, A_n \text{ are columns of } A \right\}$$

- Very easy to prove that $\text{colspan}(A)$ is a subspace of \mathbb{R}^m

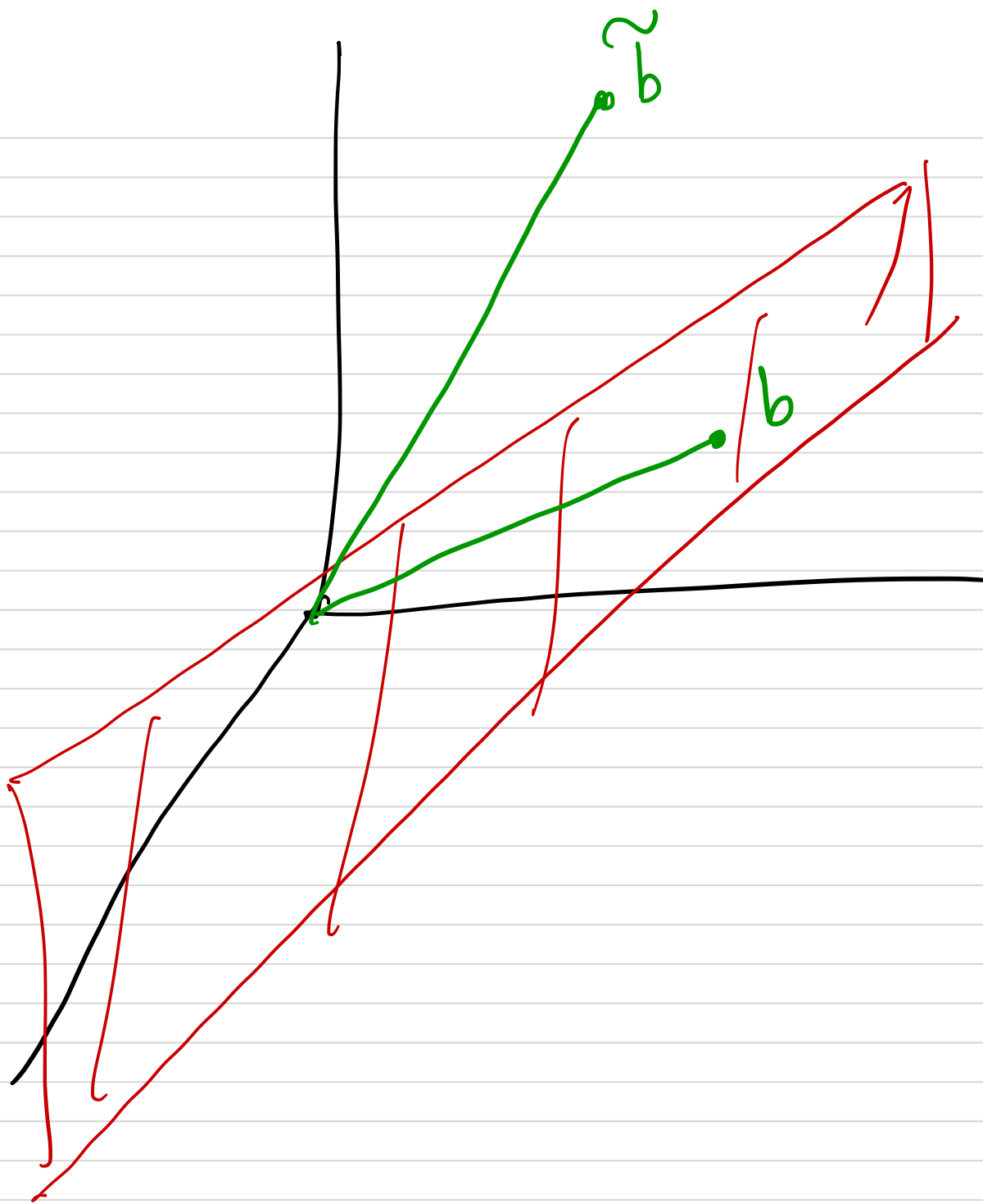
If $b \in \mathbb{R}^m$ is in the $\text{colspan}(A)$, then b can be written as a linear combination of A_i 's.

$\exists x_1, x_2, \dots, x_n \in \mathbb{R}$ s.t.

$$\sum_{i=1}^n A_i x_i = b$$

$$\Rightarrow Ax = b$$

\Rightarrow Solution exists!!



2) Uniqueness

The columns of A are in fact basis vectors for the column space of A .

Ex: Relate the geometrical concepts with algebraic constraints on $\text{rank} [A \mid b]$

Solve:

$$Ax = b$$

Direct
methods

Iterative
methods

$$\begin{bmatrix} u_{11} & u_{12} & \dots & u_{1n} \\ \star & u_{22} & \dots & u_{2n} \\ \star & & \dots & \\ \star & \star & & u_{nn} \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$$

$$Ux = b$$

$$u_{nn} x_n = b_n$$

$$x_n = \frac{1}{u_{nn}} b_n$$

$$u_{n-1,n-1} x_{n-1} + u_{n-1,n} x_n = b_{n-1}$$

$$x_{n-1} = \frac{1}{u_{n-1,n-1}} \left[b_{n-1} - u_{n-1,n} x_n \right]$$

Floating Point Operations flops
→ Floating Point Arithmetic

Similarly, $Lx = b$ where L is an lower triangular system can also be solved very easily by forward substitution.

In general, we will try to convert $Ax = b$ problem into an equivalent upper triangular or lower triangular system.

LU, cholesky, QR

Symmetric positive definite matrices.

$$A \in \mathbb{R}^{n \times n}$$

A : symmetric.

$$x^T A x > 0 \quad \forall x \in \mathbb{R}^n$$

$Ax = b$ where symmetric positive definite -

$$A = \begin{bmatrix} \square & & \\ & \square & \\ & & \square \end{bmatrix}$$

principal minors

Lecture 2 :

• Gaussian Elimination (pivoting)

$$A = \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} \\ a_{21} & a_{22} & \dots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} \end{bmatrix} ; b = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$$

$[A ; b]$

$$\boxed{R_2 - \frac{a_{21}}{a_{11}} R_1}$$

$$\begin{bmatrix} \diagdown & & & \\ & \diagdown & & \\ & & \ddots & \\ & & & \diagdown \end{bmatrix} \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_n \end{bmatrix}$$

• LU

$$A \in \mathbb{R}^{n \times n}$$

non-singular.

$$A = LDU$$

$$\begin{bmatrix} 1 & & & 0 \\ & \ddots & & \\ & & \ddots & \\ & * & & 1 \end{bmatrix} \begin{bmatrix} d_1 & & & 0 \\ & d_2 & & \\ & & \ddots & \\ 0 & & & d_n \end{bmatrix} \begin{bmatrix} 1 & & & * \\ & \ddots & & \\ & & \ddots & \\ 0 & & & 1 \end{bmatrix}$$

$$Ax = b$$

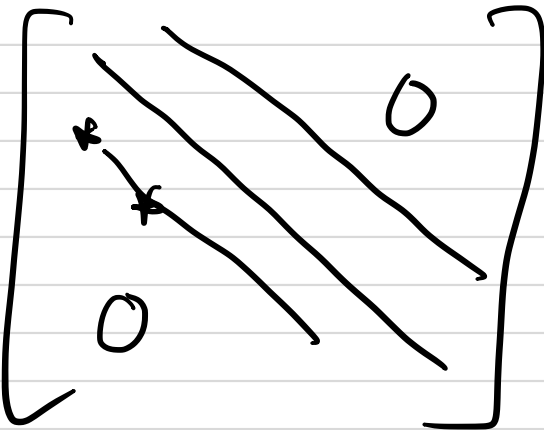
$$L[Ux] = b$$

y

$$Ly = b$$

$$Ux = y$$

$$Ax = b$$



$$A \in \mathbb{R}^{n \times n}$$

symmetric

$$A = LL^T$$

A is symmetric p.d.

$$A = LDL^T$$

where D has only positive diagonal entries.

$$A = GG^T$$

where $G = LD^{1/2}$

To solve $Ax = b$

In practice we generally solve a perturbed system.

$$(A + \Delta A)(x + \Delta x) = (b + \Delta b)$$

$$\tilde{A}\tilde{x} = \tilde{b}$$

Suppose that one can solve $\tilde{A}\tilde{x} = \tilde{b}$ exactly for \tilde{x} .

Then we declare \tilde{x} to be the solution to $Ax = b$.

Can we justify this??

Answer: Not always !! (Sensitivity analysis)

Ex: $Ax = b$

$$A = \begin{bmatrix} 1000 & 999 \\ 999 & 998 \end{bmatrix} ; b = \begin{bmatrix} 1999 \\ 1997 \end{bmatrix}$$

Note: $x = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$ is the solution.

$$\tilde{b} = \begin{bmatrix} 1998.99 \\ 1997.01 \end{bmatrix}$$

$$\tilde{b} = b + \Delta b$$

$$\Delta b = \begin{bmatrix} -0.01 \\ 0.01 \end{bmatrix} *$$

Now solve:

$$A \tilde{x} = \tilde{b}$$

$$\tilde{x} = \begin{bmatrix} 20.97 \\ -18.99 \end{bmatrix} *$$

$$A^{-1} = \begin{bmatrix} -998 & 999 \\ 999 & -1000 \end{bmatrix}$$

$$\tilde{b} = \begin{bmatrix} 1999.01 \\ 1997.01 \end{bmatrix}$$

$$\Delta b = \begin{bmatrix} 0.01 \\ 0.01 \end{bmatrix}$$

$$\tilde{x} = A^{-1} \tilde{b}$$

$$= \begin{bmatrix} -998 & 999 \\ 999 & -1000 \end{bmatrix} \begin{bmatrix} 1999.01 \\ 1997.01 \end{bmatrix}$$

$$\tilde{x} = \begin{bmatrix} -998 \times 1999.01 + 999 \times 1997.01 \\ \dots \end{bmatrix}$$

$$\tilde{x} = \begin{bmatrix} 1.01 \\ 0.99 \end{bmatrix}$$

We observe here that a "small" perturbation in b causes a "very large" perturbation in x . *

Q: Can we quantify this perturbation in x ??

Norms :

vector norm.

$$\|\cdot\| : \mathbb{R}^n \rightarrow \mathbb{R}$$

$$\text{i) } \|x\| \geq 0 \quad \forall x \in \mathbb{R}^n$$

and $\|x\| = 0$ iff $x = 0$

$$\text{ii) } \|\alpha x\| = |\alpha| \|x\|$$

$$\text{iii) } \|x+y\| \leq \|x\| + \|y\|$$

$$\text{(.) } \|x\|_2 = \sqrt{\langle x, x \rangle} = \left(\sum_{i=1}^n x_i^2 \right)^{1/2}$$

$$\text{(.) } \|x\|_p = \left(\sum_{i=1}^n |x_i|^p \right)^{1/p} \quad p \geq 1$$

is a vector norm.

$$\text{(.) } \|x\|_1 = \sum_{i=1}^n |x_i|$$

$$\text{(.) } \|x\|_\infty = \max_i |x_i|$$

Matrix norms:

$$A \in \mathbb{R}^{n \times n}$$

Consider A as a vector in \mathbb{R}^{n^2} .

$$\|A\|_F = \left(\sum_{i,j} a_{ij}^2 \right)^{1/2}$$

is in fact the 2-norm of the vector in \mathbb{R}^{n^2} .

Induced matrix norm:

$$\|A\|_2 = \max_{x \neq 0} \frac{\|Ax\|_2}{\|x\|_2}$$

$$= \max_{x \neq 0} \left\| A \frac{x}{\|x\|_2} \right\|_2$$

$$= \max_{\|y\|_2 = 1} \|Ay\|_2$$

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Lecture 4.

Least squares problem.

$$A \in \mathbb{R}^{N \times n}$$

$$b \in \mathbb{R}^N$$

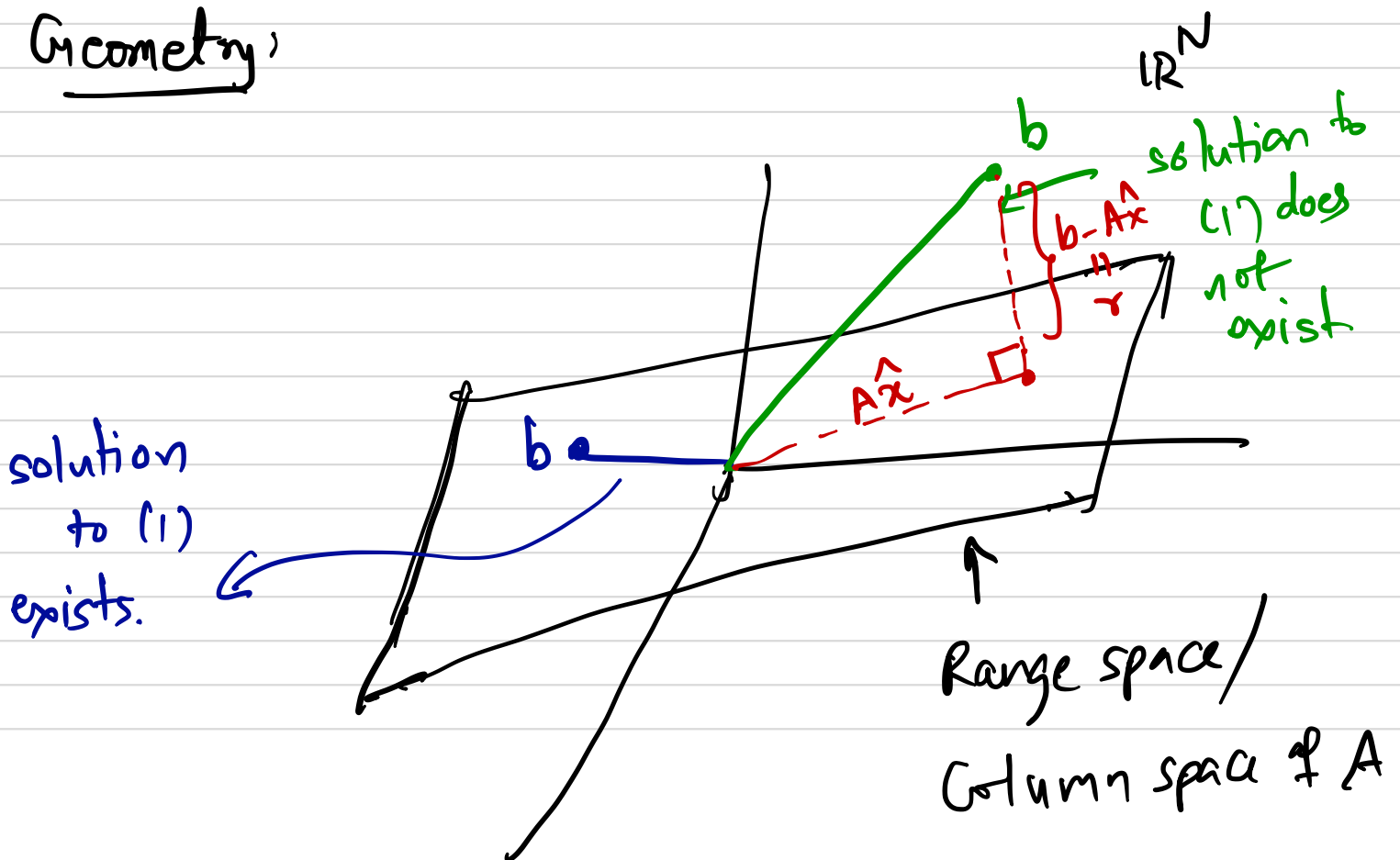
$N \gg n$. } Given

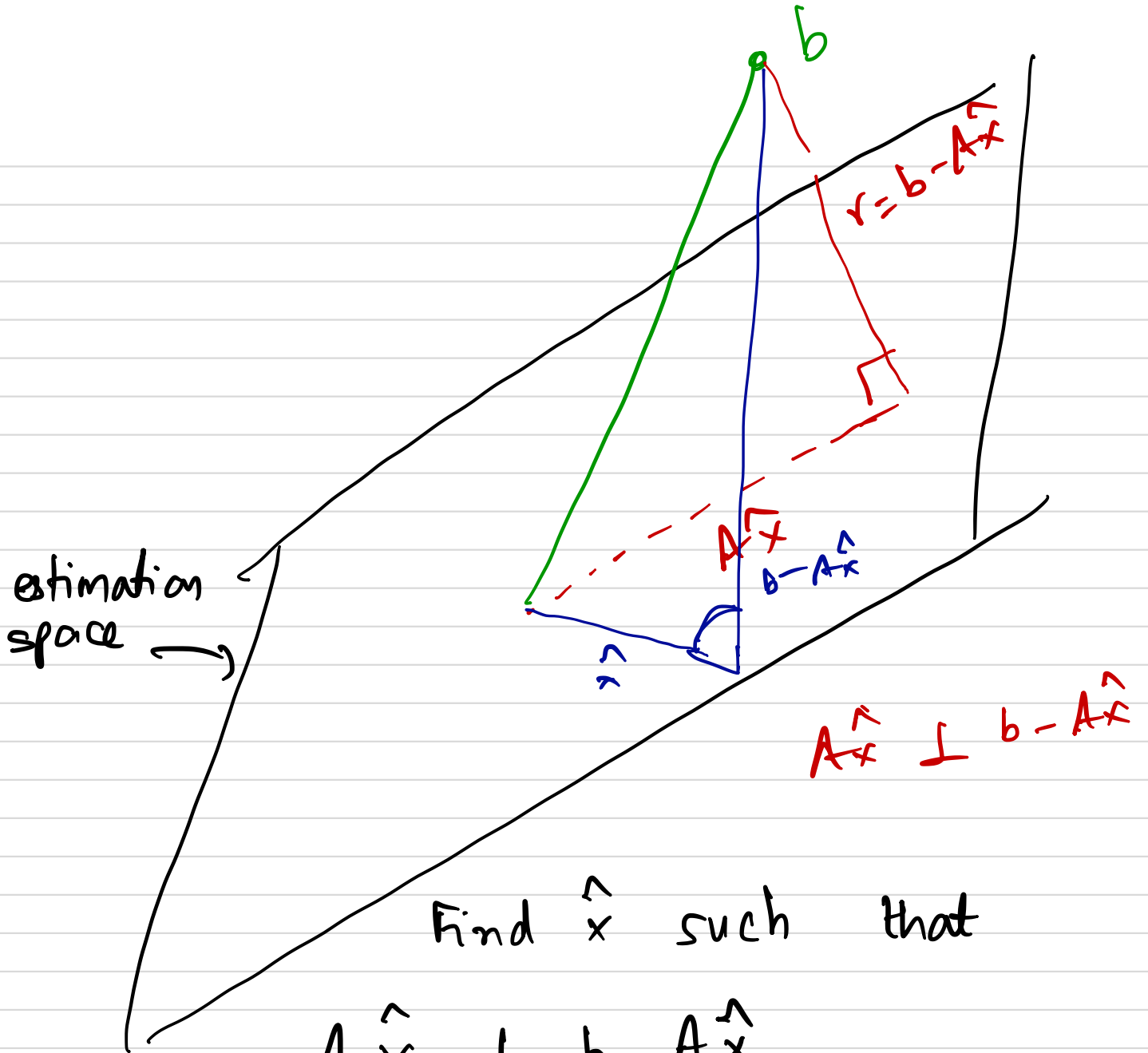
Solve: $Ax = b$

$$\begin{bmatrix} A \end{bmatrix} \begin{bmatrix} x \end{bmatrix} = \begin{bmatrix} b \end{bmatrix} \quad (1)$$

over-determined system of linear eq^s.

Geometry





$$\Rightarrow \langle A\hat{x}, b - A\hat{x} \rangle = 0$$

$$\Rightarrow \hat{x}^T A^T (b - A\hat{x}) = 0$$

$$\Rightarrow \hat{x}^T A^T b = \hat{x}^T A^T A \hat{x}$$

$$\Rightarrow \boxed{A^T b = A^T A \hat{x}} \leftarrow \text{Normal Eq.} \begin{matrix} \text{why??} \\ \swarrow \end{matrix}$$

Solving the normal eqⁿ

$$A^T A \in \mathbb{R}^{n \times n}$$

$$A^T b \in \mathbb{R}^{n \times 1}$$

$$(A^T A) \hat{x} = A^T b$$

we get the LS solution to the system $Ax = b$

In particular, if $A^T A$ is invertible,

then
$$\hat{x} = (A^T A)^{-1} A^T b$$

The matrix $(A^T A)^{-1} A^T$ is known as pseudo-inverse of A and is denoted as A^+ .

Q: Why is this solution called Least Squares (LS)??

$$\min_{x \in \mathbb{R}^n} \|b - Ax\|_2 \leftarrow$$

$$\min_{x \in \mathbb{R}^n} \|b - Ax\|_2^2 \leftarrow \textcircled{LS}$$

$$\|b - Ax\|_2^2 = \langle b - Ax, b - Ax \rangle$$

$$= (b - Ax)^T (b - Ax)$$

$$= (b^T - x^T A^T) (b - Ax)$$

$$= \underbrace{b^T b - x^T A^T b - b^T Ax}_{x^T A^T A x}$$

$$= x^T A^T A x - \underbrace{2 x^T A^T b}_{\text{const}} + b^T b$$

$$\min_{x \in \mathbb{R}^n} \underbrace{x^T A^T A x - 2 x^T A^T b + b^T b}_{\text{const } f''}$$

$$\nabla_x (\text{const } f'') = 0$$

$$2 A^T A x - 2 A^T b = 0$$

$$\Rightarrow A^T A x = A^T b$$

$$\langle b - Ax, b - Ax \rangle$$

$$x = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$$

$$b = \begin{bmatrix} \\ \\ \\ \\ \end{bmatrix}_{N \times 1} \quad Ax = \begin{bmatrix} \\ \\ \\ \\ \end{bmatrix}_{N \times 1}$$

$$r = b - Ax = \begin{bmatrix} r_1 \\ r_2 \\ \vdots \\ r_N \end{bmatrix}_{N \times 1}$$

$$\langle b - Ax, b - Ax \rangle = \langle r, r \rangle$$

$$= \sum_{i=1}^N r_i^2$$

Given A & b , to solve the LS problem, one needs to construct normal equations and then solve them simultaneously.

$$\underbrace{A^T A}_{} \hat{x} = \underbrace{A^T b}_{}$$

i) computational complexity of computing $A^T A$.

ii)
$$A = \begin{bmatrix} 1 & 1 \\ 1 & 1 \\ 0 & \epsilon \end{bmatrix}$$

$$A^T A = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & \epsilon \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 1 \\ 0 & \epsilon \end{bmatrix}$$

$$= \begin{bmatrix} 2 & 2 \\ 2 & 2 + \epsilon^2 \end{bmatrix} \quad \textcircled{A}$$

Numerical approach to solve LS problem.

$$Ax = b$$

$$A \in \mathbb{R}^{N \times n}$$

$$b \in \mathbb{R}^{N \times 1}$$

To find x s.t. $\|b - Ax\|_2^2$ is minimized.

$$\|b - Ax\|_2^2 = \langle b - Ax, b - Ax \rangle$$

$$= \langle Q(b - Ax), Q(b - Ax) \rangle$$

$$= (b - Ax)^T \underbrace{Q^T Q}_I (b - Ax)$$

If Q is such that $Q^T Q = I$

$$\|Q^T(b - Ax)\|_2^2 = \|b - Ax\|_2^2$$

$$\|Q^T b - \underbrace{Q^T A}_R x\|_2^2 = \|Q^T b - R x\|_2^2$$

$$R = \begin{bmatrix} \diagdown \end{bmatrix}$$